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# Applicability of the *q*-analogue of Zeilberger's algorithm

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#### Abstract

The applicability or terminating condition for the ordinary case of Zeilberger's algorithm was recently obtained by Abramov. For the q-analogue, the question of whether a bivariate q-hypergeometric term has a qZ-pair remains open. Le has found a solution to this problem when the given bivariate q-hypergeometric term is a rational function in certain powers of q. We solve the problem for the general case by giving a characterization of bivariate q-hypergeometric terms for which the q-analogue of Zeilberger's algorithm terminates. Moreover, we give an algorithm to determine whether a bivariate q-hypergeometric term has a qZ-pair. © 2004 Elsevier Ltd. All rights reserved.

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#### 1. Introduction

Zeilberger's algorithm (Graham et al., 1994; Petkovšek et al., 1996; Zeilberger, 1991), also known as the method of *creative telescoping*, is devised for proving hypergeometric identities of the form

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$$\sum_{k=-\infty}^{\infty} F(n,k) = f(n),$$

where F(n, k) is a bivariate hypergeometric term and f(n) is a given function (for most cases a hypergeometric term plus a constant). The algorithm can be easily adapted to the q-case, which is called the q-analogue of Zeilberger's algorithm (Böing and Koepf, 1999; Koornwinder, 1993; Paule and Riese, 1997; Wilf and Zeilberger, 1992). Let N and K be the shift operators with respect to n and k respectively, defined by

$$NT(n,k) = T(n+1,k)$$
 and  $KT(n,k) = T(n,k+1).$ 

Given a bivariate q-hypergeometric term T(n, k), the q-analogue of Zeilberger's algorithm aims to find a qZ-pair (L, G), where L is a linear difference operator with coefficients in the ring of polynomials in  $q^n$ 

$$L = a_0(q^n)N^0 + a_1(q^n)N^1 + \dots + a_r(q^n)N^n$$

and G is a bivariate q-hypergeometric term G(n, k) such that

$$LT(n,k) = (K-1)G(n,k).$$

Zeilberger's algorithm has been widely used as a powerful tool to prove hypergeometric identities. It was an open question when the algorithm terminates. This problem was solved recently by Abramov (2002, 2003). For the q-analogue of Zeilberger's algorithm, Abramov and Le (2002) found a solution to the termination problem for the case of rational functions. In this paper we provide a solution for the general q-case.

We begin with an additive decomposition of univariate q-hypergeometric terms. Using this decomposition, a univariate q-hypergeometric term T(n) can be represented as

$$T(n) = (N-1)T_1(n) + T_2(n),$$

where  $T_1(n)$  and  $T_2(n)$  are q-hypergeometric terms, and  $T_2(n)$  has the following form:

$$T_2(n) = \frac{u_1(q^n)}{u_2(q^n)} \prod_{j=n_0}^{n-1} \frac{f_1(q^j)}{f_2(q^j)},$$

where  $u_1, u_2, f_1, f_2$  are polynomials,  $n_0$  is a nonnegative integer, and for any integer m,  $u_2(x)$  and  $u_2(xq^m)$  have no common factors except for a power of x. Consequently, a bivariate q-hypergeometric term T(n, k) can be decomposed as

$$T(n,k) = (K-1)T_1(n,k) + T_2(n,k)$$
(1.1)

such that

$$T_2(n,k) = T(n,k_0)V(q^n,q^k) \prod_{j=k_0}^{k-1} F(q^n,q^j),$$

where V, F are rational functions,  $n_0$  is a nonnegative integer, and the denominator  $v_2$  of V satisfies the conditions that for any integer m,  $v_2(x, y)$  and  $v_2(x, yq^m)$  have no common factors except for a power of y. The polynomial  $v_2(x, y)$  with the above property

is called  $\varepsilon_y$ -free. We should note that the above decomposition does not solve the minimal additive decomposition problem and is not unique (see Abramov and Petkovšek (2002a) for a precise definition). However, for the purpose of constructing a *qZ*-pair, it turns out that one may choose any decomposition.

Then we consider the structure of bivariate q-hypergeometric terms. The structure of ordinary hypergeometric terms has been studied by Ore (1930), Sato et al. (1990), Gel'fand et al. (1992), Abramov and Petkovšek (2002b) and Hou (2004). To a large extent, the q-case is analogous to the ordinary case. For each bivariate q-hypergeometric term, we associate it with a normal representation (q-NR) which consists of four polynomials r, s, u, v. Based on the properties of the representation, we may give a definition of q-proper hypergeometric terms and prove that under the condition that v is  $\varepsilon_y$ -free, a bivariate q-hypergeometric term has a qZ-pair if and only if it is a q-proper term. Applying the decomposition (1.1), we deduce that for any bivariate q-hypergeometric term T, it has a qZ-pair if and only if  $T_2$  is q-proper.

We conclude with some examples.

#### 2. $\varepsilon$ -free decomposition

Throughout the paper, we let  $\mathbb{Z}$ ,  $\mathbb{Z}^+$  and  $\mathbb{N}$  denote the set of integers, positive integers and nonnegative integers, respectively. For integers (or polynomials) *a*, *b*, we denote by gcd(*a*, *b*) the (monic) greatest common divisor of *a* and *b*. We also write  $a \perp b$  to indicate that *a* and *b* are relatively prime, i.e., gcd(*a*, *b*) = 1.

Let  $\mathbb{F}$  be a field of characteristic zero,  $q \in \mathbb{F}$  a nonzero element which is not a root of unity, and x transcendental over  $\mathbb{F}$ . Denote by  $\boldsymbol{\varepsilon}$  the unique automorphism of  $\mathbb{F}(x)$  which fixes  $\mathbb{F}$  and satisfies  $\boldsymbol{\varepsilon} x = qx$ . Then  $\mathbb{F}(x)$  together with the *q*-shift operator  $\boldsymbol{\varepsilon}$  is a difference field (Cohn, 1965). Let *r* and *s* be two polynomials. We say that r/s is  $\boldsymbol{\varepsilon}$ -reduced if  $r \perp \boldsymbol{\varepsilon}^h s$  for all  $h \in \mathbb{Z}$ .

To be more specific, the rational functions involved in the *q*-hypergeometric terms (see Definition 2.4) are rational functions of  $q^n$ . However, for a rational function  $R \in \mathbb{F}(x)$  and a nonnegative integer  $n_0$ , we have

$$N R(q^n) = R(q^{n+1}) = \varepsilon R(q^n)$$
 and  $R(q^n) = 0 \forall n \ge n_0 \Leftrightarrow R(x) = 0$ .

Therefore, there is a natural one-to-one correspondence between the set of rational functions of  $q^n$  together with the shift operator N and the field  $\mathbb{F}(x)$  together with the q-shift operator  $\boldsymbol{\varepsilon}$ . In this paper, we adopt the notation of  $\mathbb{F}(x)$  as in the work of Abramov et al. (1998).

The concept of rational normal forms introduced by Abramov and Petkovšek (2002a) can be extended to the q-case.

**Definition 2.1.** Let  $R \in \mathbb{F}(x)$  be a rational function. If polynomials  $r, s, u, v \in \mathbb{F}[x]$  satisfy

(i)  $R = \frac{r}{s} \cdot \frac{\varepsilon(u/v)}{(u/v)}$ , where  $u \perp v$  and u, v have no factor x, (ii) r/s is  $\varepsilon$ -reduced,

then (r, s, u, v) is called a *q*-rational normal form (q-RNF) of *R*.

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Recall that a monic polynomial that has no factor x is called a q-monic polynomial by Abramov et al. (1998). The following factorization theorem was given in Abramov et al. (1998).

**Theorem 2.2.** Let  $R \in \mathbb{F}(x) \setminus \{0\}$ . Then there exist  $z \in \mathbb{F}$  and monic polynomials  $a, b, c \in \mathbb{F}[x]$  such that

$$R(x) = z \frac{a(x)}{b(x)} \frac{c(qx)}{c(x)},$$
  

$$gcd(a(x), b(q^n x)) = 1, \quad for \ all \ n \in \mathbb{N},$$
  

$$gcd(a(x), c(x)) = gcd(b(x), c(qx)) = 1 \quad and \quad c(0) \neq 0.$$
(2.1)

We call (az, b, c) a *q*-Gosper form (q-GF) of *R*.

**Theorem 2.3.** *Every rational function*  $R \in \mathbb{F}(x)$  *has a q-RNF.* 

**Proof.** It is clear that (0, 1, 1, 1) is a *q*-RNF of 0. For  $R \neq 0$ , by Theorem 2.2, there exists a *q*-GF (az, b, c) of *R*. Applying Theorem 2.2 again to b(x)/a(x), we get a *q*-GF (r, s, d). From the construction given in Abramov et al. (1998), we have  $r \mid b$  and  $s \mid a$ . Hence  $s(x) \perp r(xq^n)$  for any  $n \in \mathbb{N}$  because (az, b, c) is a *q*-GF. Since (r, s, d) is also a *q*-GF, we have  $r(x) \perp s(xq^n)$  for any  $n \in \mathbb{N}$ . Thus s/r is  $\varepsilon$ -reduced and  $(zs, r, c/\gcd(c, d), d/\gcd(c, d))$  is a *q*-RNF of *R*.  $\Box$ 

The above proof provides an algorithm to generate a q-RNF of R.

#### Algorithm q-RNF

if R = 0 then
 return (0, 1, 1, 1);
else
 compute 'q-GF' of R, we get (a, b, c);
 compute 'q-GF' of b/a, we get (r, s, d);
 return (s, r, c/ gcd(c, d), d/ gcd(c, d)).

We now come to the q-multiplicative representation of a general q-hypergeometric term. This is the starting point of the  $\varepsilon$ -free decomposition algorithm.

**Definition 2.4.** Suppose T(n) is a function from  $\mathbb{N}$  to  $\mathbb{F}$ . If there exist a nonnegative integer  $n_0$  and a nonzero rational function  $R(x) \in \mathbb{F}(x)$  such that  $T(n + 1) = R(q^n)T(n)$  for all  $n \ge n_0$ , then we call T(n) a (univariate) *q*-hypergeometric term.

Suppose (r, s, u, v) is a *q*-RNF of a rational function *R*. Then the corresponding *q*-hypergeometric term T(n) satisfies

$$T(n) = T(n_0) \prod_{j=n_0}^{n-1} R(q^j) = \frac{T(n_0)}{u(q^{n_0})/v(q^{n_0})} \cdot \frac{u(q^n)}{v(q^n)} \prod_{j=n_0}^{n-1} \frac{r(q^j)}{s(q^j)}, \quad \forall n \ge n_0.$$

This leads to the following definition.

$$T(n) = U(q^n) \prod_{j=n_0}^{n-1} D(q^j), \qquad \forall n \ge n_0.$$

Then we call  $(D, U, n_0)$  a q-multiplicative representation (q-MR) of T.

Let  $\Delta = N - 1$  be the difference operator with respect to *n*. The following lemma can be easily verified.

**Lemma 2.6.** Let T and  $T_1$  be two q-hypergeometric terms with q-MRs  $(D, U, n_0)$  and  $(D, U_1, n_0)$ , respectively. Suppose that

$$T_2 = T - \Delta T_1$$
 and  $U_2 = U - D \cdot \boldsymbol{\varepsilon} U_1 + U_1$ .

Then  $(D, U_2, n_0)$  is a q-MR of  $T_2$ .

For  $u, v \in \mathbb{F}[x]$ , let  $\mathcal{R}$  be the set of all nonnegative integers h such that there exists an irreducible polynomial  $p(x) \neq x$  satisfying p(x) | u(x) and  $p(x) | v(q^h x)$ . Define qdis(u, v) to be  $max\{h \in \mathcal{R}\}$  or -1 if  $\mathcal{R}$  is empty. Note that  $\mathcal{R}$  is a finite set, and "qdis" is well defined. If qdis(v, v) = 0, we say that v is  $\varepsilon$ -free.

Given a *q*-hypergeometric term *T* with a *q*-MR  $(D, U, n_0)$ . Usually the denominator *u* of *U* is not  $\boldsymbol{\varepsilon}$ -free. However, translating the decomposition algorithm of Abramov and Petkovšek (2002a) into the *q*-case, we have the following  $\boldsymbol{\varepsilon}$ -free decomposition algorithm "*q*-decomp", which decomposes *T* into  $\Delta T_1 + T_2$  such that  $T_2$  has a *q*-MR  $(F, V, n_0)$  where the denominator of *V* is  $\boldsymbol{\varepsilon}$ -free.

#### Algorithm q-decomp

```
Input: (D, U, n_0)
                                 Output: U_1, F, V \in \mathbb{F}(x)
    d_1 := \operatorname{numer}(D); d_2 := \operatorname{denom}(D);
    U_1 := 0; U_2 := U; u_2 := \text{denom}(U);
    N := \operatorname{qdis}(u_2, u_2);
    for h := N down to 1 do
        v_2 := u_2 / \gcd(u_2, d_2);
        s(x) := \gcd(v_2(x), v_2(q^{-h}x));
        (\tilde{s}, \tilde{u}_2) := \operatorname{pump}(s, u_2);
        write U_2 = a/\tilde{u}_2 + b/\tilde{s} where a, b \in \mathbb{F}[x];
        U'_{1} := -b/\tilde{s};
        U_1 := U_1 + U'_1; U_2 := U_2 - D \cdot \varepsilon U'_1 + U'_1;
        u_2 := \operatorname{denom}(U_2);
    f_1 := d_1; f_2 := d_2; v_1 := \operatorname{numer}(U_2); v_2 := \operatorname{denom}(U_2);
    w := \gcd(d_2, v_2);
    v_2 := v_2/w; f_2 := \varepsilon w f_2/w;
    F := f_1/f_2; V := (1/w(q^{n_0})) \cdot v_1/v_2;
    return (U_1, F, V).
```

The procedure "pump" is the same as in the ordinary case.

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#### Algorithm pump

Input:  $f, g \in \mathbb{F}[x]$  Output:  $\tilde{f}, \tilde{g} \in \mathbb{F}[x]$ 

$$\begin{split} \tilde{f} &:= f; \, \tilde{g} := g/f; \\ \text{repeat} \\ d &:= \gcd(\tilde{f}, \, \tilde{g}); \quad \tilde{f} := \tilde{f}d; \, \tilde{g} := \tilde{g}/d; \\ \text{until } \deg d &= 0; \\ \text{return } (\tilde{f}, \, \tilde{g}). \end{split}$$

The following theorem shows that the  $\boldsymbol{\varepsilon}$ -free algorithm generates the desired decomposition.

**Theorem 2.7.** Let T be a q-hypergeometric term with a q-MR  $(D, U, n_0)$  and  $U_1, F, V$  be given by the algorithm q-decomp. Then there exist q-hypergeometric terms  $T_1$  and  $T_2$  such that

- (1)  $T = \Delta T_1 + T_2$ .
- (2)  $T_1$  has a q-MR  $(D, U_1, n_0)$  and  $T_2$  has a q-MR  $(F, V, n_0)$ .
- (3) The denominator of V is  $\boldsymbol{\varepsilon}$ -free.

Furthermore, if D is  $\varepsilon$ -reduced, so is F.

**Proof.** Let  $u_0$  be the denominator of U. We first use induction to show that after iterating the loop of h in the algorithm i times, the denominator  $u_2$  of  $U_2$  satisfies:

(a) qdis(v<sub>2</sub>, v<sub>2</sub>) ≤ N − i,
(b) u<sub>2</sub>(q<sup>n</sup>) has no zeros for all n ≥ n<sub>0</sub>,

where  $v_2 = u_2 / \operatorname{gcd}(u_2, d_2)$ , and  $d_2$  is the denominator of *D*.

The case for i = 0 is trivial. Assume that the assertion holds for i - 1. Let  $u_2$  and  $u'_2$  be the denominator of  $U_2$  after i - 1 and i iterations, respectively. Set h = N - (i - 1) > 0 and  $w_2 = \text{gcd}(u_2, d_2)$ . From the algorithm q-decomp we have

$$v_2 = u_2/w_2$$
 and  $s = \gcd(v_2(x), v_2(q^{-n}x))$ 

Suppose the prime decomposition of s is  $p_1^{\alpha_1} \cdots p_r^{\alpha_r}$  and  $v_2 = p_1^{\beta_1} \cdots p_r^{\beta_r} v', w_2 = p_1^{\gamma_1} \cdots p_r^{\gamma_r} w'$  where  $v' \perp s, w' \perp s$ . Then the algorithm "pump" enables us to decompose  $u_2$  as  $p_1^{\beta_1+\gamma_1} \cdots p_r^{\beta_r+\gamma_r} \cdot (v'w')$ . That is,  $\tilde{s} = p_1^{\beta_1+\gamma_1} \cdots p_r^{\beta_r+\gamma_r}$  and  $\tilde{u}_2 = v'w'$ . Since

$$U_2 = \frac{a}{\tilde{u}_2} + \frac{d_1}{d_2} \cdot \boldsymbol{\varepsilon}\left(\frac{b}{\tilde{s}}\right),$$

it follows that  $u'_2$  divides the least common multiple of  $\tilde{u}_2$  and  $d_2 \epsilon \tilde{s}$ . Hence we have that  $u'_2$  divides  $v'd_2 \cdot \epsilon \tilde{s}$ . Let  $v'' = v' \cdot \epsilon \tilde{s}$ . Assume that there exist an integer  $m \ge h$  and an irreducible polynomial  $p(x) \ne x$  such that p | v'' and  $p | \epsilon^m v''$ . We may encounter four cases:

•  $p \mid v'$  and  $p \mid \boldsymbol{\varepsilon}^m v'$ .

From  $v' | v_2$  and  $qdis(v_2, v_2) \le h$ , it follows that m = h. Therefore,  $\varepsilon^{-h}p | \varepsilon^{-h}v_2$  and  $\varepsilon^{-h}p | v_2$ . Consequently, we have  $\varepsilon^{-h}p | s$ , which contradicts  $v' \perp s$ .

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•  $p \mid v' \text{ and } p \mid \boldsymbol{\varepsilon}^{m+1} \tilde{s}$ .

Since *s* and  $\tilde{s}$  have the same prime factors, we have  $p | \boldsymbol{\varepsilon}^{m+1} s$ , implying that  $p | \boldsymbol{\varepsilon}^{m+1} v_2$ . On the other hand, we have  $p | v_2$ , which contradicts  $qdis(v_2, v_2) \le h$ .

- $p | \boldsymbol{\varepsilon} \tilde{s}$  and  $p | \boldsymbol{\varepsilon}^m v'$ . In this situation, we have  $\boldsymbol{\varepsilon}^{-1} p | \tilde{s}$ , which implies that  $\boldsymbol{\varepsilon}^{-1} p | \boldsymbol{\varepsilon}^{-h} v_2$ , or equivalently,  $\boldsymbol{\varepsilon}^{h-1} p | v_2$ . On the other hand,  $\boldsymbol{\varepsilon}^{h-1} p | \boldsymbol{\varepsilon}^{m+h-1} v_2$ . Since  $\operatorname{qdis}(v_2, v_2) \leq h$ , we get  $m+h-1 \leq h$ , and hence m = 1. Now we have  $p | \boldsymbol{\varepsilon} s$  and  $p | \boldsymbol{\varepsilon} v'$ , which contradicts  $v' \perp s$ .
- $p | \boldsymbol{\varepsilon}\tilde{s}$  and  $p | \boldsymbol{\varepsilon}^{m+1}\tilde{s}$ . Similarly, we have  $\boldsymbol{\varepsilon}^{-1}p | s$  and hence  $\boldsymbol{\varepsilon}^{-1}p | \boldsymbol{\varepsilon}^{-h}v_2$ , i.e.,  $\boldsymbol{\varepsilon}^{h-1}p | v_2$ . However, we have  $\boldsymbol{\varepsilon}^{h-1}p | \boldsymbol{\varepsilon}^{m+h}v_2$ . Thus, we obtain  $m+h \leq h$ , which is also a contradiction.

In summary, we may conclude that  $\operatorname{qdis}(v'', v'') \leq h - 1$ . Because  $u'_2$  divides  $v'' \cdot d_2$ , there exist  $\bar{v} \mid v''$  and  $\bar{w} \mid d_2$  such that  $u'_2 = \bar{v}\bar{w}$ . Let  $v'_2 = u'_2/\operatorname{gcd}(u'_2, d_2)$ . From  $\bar{w} \mid \operatorname{gcd}(u'_2, d_2)$ , it follows that  $v'_2 \mid \bar{v}$ . So we get  $\operatorname{qdis}(v'_2, v'_2) \leq h - 1 = N - i$ . Thus, we have proved (a). Since  $u'_2 \mid u_2 \cdot \epsilon u_2 \cdot d_2$ , (b) immediately follows from the induction hypothesis.

On the other hand, since  $\tilde{s} \mid u_2$ , (b) implies that  $U_1(q^n)$  has no poles for all  $n \ge n_0$ . Let

$$T_1(n) = U_1(q^n) \prod_{j=n_0}^{n-1} D(q^j)$$
 and  $T_2(n) = U_2(q^n) \prod_{j=n_0}^{n-1} D(q^j).$  (2.2)

Noting that  $U_2 = U - D\varepsilon U_1 + U_1$ , by Lemma 2.6, we obtain  $T = \Delta T_1 + T_2$ . Because  $w \mid d_2$  and  $d_2(q^n) \neq 0$  for all  $n \ge n_0$ , we can write  $T_2(n)$  as

$$T_2(n) = \frac{1}{w(q^{n_0})} U_2(q^n) w(q^n) \prod_{j=n_0}^{n-1} D(q^j) \frac{w(q^j)}{w(q^{j+1})} = V(q^n) \prod_{j=n_0}^{n-1} F(q^j).$$

Let v be the denominator of V. Then (a) implies qdis(v, v) = 0; that is, v is  $\varepsilon$ -free. Finally, notice that  $f_1 = d_1$  and  $f_2 = \varepsilon w \cdot (d_2/w)$ , where  $w \mid d_2$ . Therefore, F is

 $\boldsymbol{\varepsilon}$ -reduced provided that *D* is  $\boldsymbol{\varepsilon}$ -reduced. This completes the proof.  $\Box$ 

#### 3. Bivariate q-hypergeometric terms

We begin this section with the definition of bivariate *q*-hypergeometric terms.

**Definition 3.1.** Suppose T(n, k) is a function from  $\mathbb{N}^2$  to  $\mathbb{F}$ . If there exist rational functions  $R_1(x, y), R_2(x, y) \in \mathbb{F}(x, y)$  and  $n_0 \in \mathbb{N}$  such that

$$T(n+1,k) = R_1(q^n, q^k)T(n,k)$$
 and  $T(n, k+1) = R_2(q^n, q^k)T(n,k),$ 

for all  $n, k \ge n_0$ , then we call T(n, k) a bivariate *q*-hypergeometric term.

Without loss of generality, from now on we may assume that  $n_0 = 0$  and that  $R_1(q^n, q^k), R_2(q^n, q^k)$  have neither zeros nor poles for all  $n, k \ge 0$ .

Denote by  $\boldsymbol{\varepsilon}_x$  and  $\boldsymbol{\varepsilon}_y$  the shift operators on  $\mathbb{F}(x, y)$  defined by  $\boldsymbol{\varepsilon}_x x = qx, \boldsymbol{\varepsilon}_x|_{\mathbb{F}(y)} = \mathrm{id}$ (the identity map) and  $\boldsymbol{\varepsilon}_y y = qy, \boldsymbol{\varepsilon}_y|_{\mathbb{F}(x)} = \mathrm{id}$ , respectively. The idea of *q*-RNFs can be easily adopted to the bivariate case by taking  $\mathbb{F}(y)$  as the ground field. Let R(x, y) be

a rational function of x and y; its q-rational normal form (q-RNF with respect to  $\varepsilon_x$ ) is represented by (r, s, u, v) as in the univariate case. By using the ground field  $\mathbb{F}(x)$ , we may find a q-RNF of R(x, y) with respect to  $\varepsilon_y$ .

Let T(n, k) be a bivariate q-hypergeometric term. By definition, there exists a rational function R such that

$$T(n+1,k)/T(n,k) = R(q^n, q^k).$$

Suppose (r, s, u, v) is a *q*-RNF of *R* with respect to  $\varepsilon_x$ . We call (r, s, u, v) a *q*-normal representation (q-NR) of T(n, k) with respect to the shift operator *N*. Similarly, we can define the *q*-NR of T(n, k) with respect to the shift operator *K*.

We next give a characterization of the polynomials involved in the q-NR of bivariate q-hypergeometric terms.

**Theorem 3.2.** Let T(n,k) be a bivariate q-hypergeometric term that has a q-NR (r, s, u, v) with respect to N. Then r and s are products of polynomials having the form

$$(x^c y^d) \cdot \prod_{l=1}^a p(q^{w_l} x^a y^b),$$

where p is a Laurent polynomial of one variable,  $a \in \mathbb{Z}^+$ , b, c, d,  $w_l \in \mathbb{Z}$ ,  $a \perp b$ , and  $w_i \neq w_j \pmod{a}$ ,  $\forall i \neq j$ .

Similarly, suppose (r, s, u, v) is a q-NR of T with respect to K. Then r and s are products of polynomials having the form

$$(x^c y^d) \cdot \prod_{l=1}^a p(q^{w_l} x^b y^a)$$

under the same conditions.

**Sketch of the proof.** The proof of the ordinary case (Hou, 2004, Theorem 3.4) can be carried over to the *q*-case except that we need to consider the characterization of polynomials f(x, y) such that  $f(q^a x, q^b y) = Cf(x, y)$  for certain integers *a*, *b* and  $C \in \mathbb{F}$ .  $\Box$ 

Consequently, we have

**Corollary 3.3.** Let T(n,k) be a bivariate q-hypergeometric term that has a q-NR (r, s, u, v) with respect to N (or K respectively). Then we have

$$T(n,k) = C \cdot \frac{u(q^{n}, q^{k})}{v(q^{n}, q^{k})} \cdot \frac{\prod_{l=1}^{uu} \prod_{j=0}^{a_{l}n+b_{l}k+c_{l}} f_{l}(q^{j})}{\prod_{l=1}^{vv} \prod_{j=0}^{a'_{l}n+b'_{l}k+c'_{l}} \prod_{g_{l}(q^{j})},$$

where  $C \in \mathbb{F}$ ,  $uu, vv \in \mathbb{N}$ ,  $a_l, b_l, c_l, a'_l, b'_l, c'_l \in \mathbb{Z}$  and  $f_l, g_l$  are polynomials.

Corollary 3.3 enables us to give the following definition of q-proper hypergeometric terms.

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**Definition 3.4.** A polynomial  $f \in \mathbb{F}[x, y]$  is said to be *q*-proper if, for each of its irreducible factors  $p(x, y) \in \mathbb{F}[x, y]$ , there exist  $a, b \in \mathbb{Z}$ , not both zeros, such that  $p(x, y)|p(q^ax, q^by)$ . A bivariate *q*-hypergeometric term *T* is said to be *q*-proper if *v* is a *q*-proper polynomial, where (r, s, u, v) is a *q*-NR of *T* with respect to *N* or *K*.

Suppose that T is a bivariate q-hypergeometric term that has a q-NR (r, s, u, v) with respect to N (or K). Theorem 3.2 guarantees that r and s are both q-proper polynomials.

As in the case of ordinary bivariate hypergeometric terms (Hou, 2004, Theorem 4.2), we have an analogous "fundamental theorem" for the q-case.

**Theorem 3.5.** Let T(n, k) be a bivariate q-hypergeometric term. Then T is q-proper if and only if there exist polynomials  $a_{ij}(x) \in \mathbb{F}[x]$ , not all zero, such that

$$\sum_{0 \le i \le I, \ 0 \le j \le J} a_{ij}(q^n) T(n+i,k+j) = 0 \quad \forall n,k \ge 0.$$

Based on an analogous argument for the ordinary case as in Petkovšek et al. (1996, Theorem 6.2.1), we get

**Corollary 3.6.** Any *q*-proper hypergeometric term has a *q*Z-pair.

#### 4. The existence of *qZ*-pairs

In this section, we obtain a necessary and sufficient condition for the existence of qZ-pairs for any bivariate q-hypergeometric term based on its q-NR with respect to K. From Theorem 3.2, we have

**Corollary 4.1.** Let T(n,k) be a bivariate q-hypergeometric term that has a q-NR (r, s, u, v) with respect to K. Then there exist polynomials  $f_i(x), g_i(x) \in \mathbb{F}[x]$  and  $a_i, a'_i, b_i, b'_i \in \mathbb{Z}$  such that

$$\prod_{j=0}^{k-1} \left( \frac{r(q^{n+1}, q^j)}{r(q^n, q^j)} \cdot \frac{s(q^n, q^j)}{s(q^{n+1}, q^j)} \right) = \prod_{i=1}^{\ell} \frac{f_i(q^{a_ik+b_in})}{g_i(q^{a'_ik+b'_in})}.$$

We need to consider the following ratio:

$$\frac{T(n+i,k)}{T(n,k)} = \frac{T(n+i,0)}{T(n,0)} \prod_{j=0}^{k-1} \left\{ \frac{T(n+i,j+1)}{T(n+i,j)} \frac{T(n,j)}{T(n,j+1)} \right\},\,$$

which can be rewritten as

$$\frac{T(n+i,k)}{T(n,k)} = \prod_{l=0}^{i-1} \prod_{j=0}^{k-1} \left\{ \frac{r(q^{n+l+1},q^j)}{r(q^{n+l},q^j)} \frac{s(q^{n+l},q^j)}{s(q^{n+l+1},q^j)} \right\} \prod_{l=0}^{i-1} \frac{T(n+l+1,0)}{T(n+l,0)} \\
\times \frac{u(q^{n+i},q^k)}{u(q^{n+i},q^0)} \frac{u(q^n,q^0)}{u(q^n,q^k)} \frac{v(q^{n+i},q^0)}{v(q^{n+i},q^k)} \frac{v(q^n,q^k)}{v(q^n,q^0)}.$$
(4.1)

From Corollary 4.1 we get the following expression.

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**Lemma 4.2.** Let T(n, k) be a bivariate q-hypergeometric term that has a q-NR (r, s, u, v) with respect to K. Then for each  $i \ge 0$ , there exist q-proper polynomials  $w_1^{(i)}(x, y)$  and  $w_2^{(i)}(x, y)$  such that

$$\frac{T(n+i,k)}{T(n,k)} = \frac{u(q^{n+i},q^k)}{v(q^{n+i},q^k)} \cdot \frac{v(q^n,q^k)}{u(q^n,q^k)} \cdot \frac{w_1^{(i)}(q^n,q^k)}{w_2^{(i)}(q^n,q^k)}, \quad \forall n,k \ge 0.$$
(4.2)

An  $\boldsymbol{\varepsilon}_{v}$ -free polynomial that is not q-proper has a special factor.

**Lemma 4.3.** Let  $f \in \mathbb{F}[x, y]$  be a non-q-proper and  $\varepsilon_y$ -free polynomial. Then there exists an irreducible factor p of f such that

$$p(x, y) \perp p(q^{i}x, q^{j}y), \quad \forall (i, j) \in \mathbb{Z}^{2} \setminus \{(0, 0)\},$$
  

$$p(x, y) \perp f(q^{i}x, q^{j}y), \quad \forall (i, j) \in (\mathbb{N} \times \mathbb{Z}) \setminus \{(0, 0)\}.$$
(4.3)

**Proof.** Since f(x, y) is non-*q*-proper, by definition it has an irreducible factor  $p_1(x, y)$  such that  $p_1(x, y) \perp p_1(q^i x, q^j y), \forall (i, j) \in \mathbb{Z}^2 \setminus \{(0, 0)\}.$ 

We may factor f(x, y) as

$$f(x, y) = p_1^{\alpha_1}(q^{a_1}x, q^{b_1}y) \cdots p_1^{\alpha_r}(q^{a_r}x, q^{b_r}y) f_1(x, y),$$

where  $(a_i, b_i) \in \mathbb{Z}^2$  are distinct pairs,  $\alpha_i \in \mathbb{Z}^+$ , and  $p_1(q^i x, q^j y) \perp f_1(x, y)$  for all  $i, j \in \mathbb{Z}$ . Since f(x, y) is  $\boldsymbol{e}_y$ -free, it follows that  $a_i \neq a_j$  as long as  $i \neq j$ . Without loss of generality, we may assume that  $a_1 < a_2 < \cdots < a_r$ . Thus,  $p(x, y) = p_1(q^{a_1}x, q^{b_1}y)$  satisfies the condition (4.3).  $\Box$ 

We are now ready to give a criterion for the existence of qZ-pairs.

**Theorem 4.4.** Let T(n,k) be a bivariate q-hypergeometric term that has a q-NR (r, s, u, v) with respect to K such that v is  $\boldsymbol{e}_y$ -free. Then T(n,k) has a qZ-pair if and only if v is a q-proper polynomial.

**Proof.** Because of Corollary 3.6, it suffices to show that if T(n, k) has a qZ-pair, then it is q-proper. To this end, we assume that T(n, k) is a bivariate q-hypergeometric term. Moreover, we assume that T(n, k) is not q-proper, but it has a qZ-pair. We proceed to find a contradiction.

Clearly, for a difference operator  $L \in \mathbb{F}[q^n, N]$ , we have

$$(N \cdot L)T(n,k) = (K-1)G(n,k) \Longleftrightarrow LT(n,k) = (K-1)G(n-1,k).$$

Therefore, we may assume that T(n, k) has a qZ-pair (L, G) of the form

$$L = \sum_{i=0}^{I} a_i(q^n) N^i,$$

where  $a_i(q^n)$  are polynomials in  $q^n$  and  $a_0 \neq 0$ . Since LT/T and (K-1)G/G are both rational functions of  $q^n$  and  $q^k$ , we may assume that

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$$G(n,k) = \frac{f(q^n, q^k)}{g(q^n, q^k)}T(n,k),$$

where  $f, g \in \mathbb{F}[x, y]$  are two relatively prime polynomials.

By the definition of qZ-pairs, we have

$$\sum_{i=0}^{I} a_i(q^n) \frac{T(n+i,k)}{T(n,k)} = \frac{f(q^n, q^{k+1})}{g(q^n, q^{k+1})} \frac{T(n,k+1)}{T(n,k)} - \frac{f(q^n, q^k)}{g(q^n, q^k)}.$$
(4.4)

Substituting (4.2) into (4.4), we obtain

$$\sum_{i=0}^{I} a_i(x) \frac{u(q^i x, y)}{v(q^i x, y)} \frac{w_1^{(i)}(x, y)}{w_2^{(i)}(x, y)} = \frac{f(x, qy) r(x, y)}{g(x, qy)} \frac{r(x, y) u(x, qy)}{v(x, qy)} - \frac{f(x, y) u(x, y)}{g(x, y)} \frac{u(x, y)}{v(x, y)}.$$
(4.5)

Let  $u_1 = u/\operatorname{gcd}(u, g)$ ,  $g_1 = g/\operatorname{gcd}(u, g)$ . Multiplying

$$g_1(x,qy)g_1(x,y)v(x,qy)s(x,y)\prod_{j=0}^{I}v(q^jx,y)w_2^{(j)}(x,y)$$

to both sides of (4.5), we arrive at

$$g_{1}(x,qy)g_{1}(x,y)v(x,qy)s(x,y)$$

$$\times \sum_{i=0}^{I} a_{i}(x)u(q^{i}x,y)w_{1}^{(i)}(x,y)\prod_{j\neq i}v(q^{j}x,y)w_{2}^{(j)}(x,y)$$

$$= f(x,qy)r(x,y)u_{1}(x,qy)g_{1}(x,y)\prod_{j=0}^{I}v(q^{j}x,y)w_{2}^{(j)}(x,y)$$

$$- f(x,y)u_{1}(x,y)g_{1}(x,qy)v(x,qy)s(x,y)w_{2}^{(0)}(x,y)$$

$$\times \prod_{j=1}^{I}v(q^{j}x,y)w_{2}^{(j)}(x,y).$$
(4.6)

Since T(n, k) is not q-proper, from Lemma 4.3 it follows that there exists an irreducible factor p of v satisfying the condition (4.3). Noting that p(x, y) divides each term of the left-hand side of (4.6) except for the first term, we obtain that p(x, y) divides

$$g_1(x, qy)v(x, qy)s(x, y) \prod_{j=1}^{I} v(q^j x, y)w_2^{(j)}(x, y) \\ \times (g_1(x, y)a_0(x)u(x, y)w_1^{(0)}(x, y) + f(x, y)u_1(x, y)w_2^{(0)}(x, y)).$$

From (4.3) it follows that

$$p(x, y) \perp v(x, qy) \prod_{j=1}^{I} v(q^{j}x, y).$$

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Since s and  $w_2^{(j)}$  are q-proper, they are also relatively prime to p. This implies that p(x, y) divides

$$g_1(x, qy) \Big( g_1(x, y) a_0(x) u(x, y) w_1^{(0)}(x, y) + f(x, y) u_1(x, y) w_2^{(0)}(x, y) \Big).$$
(4.7)

Similarly, since p(x, qy) divides both sides of (4.6) and  $u \perp v$ , we have

$$p(x, qy) | f(x, qy)g_1(x, y).$$
(4.8)

Case 1. Suppose p(x, qy) | f(x, qy). Since p(x, y) divides (4.7), it follows that

 $p(x, y) | g_1(x, qy)g_1(x, y)a_0(x)u(x, y)w_1^{(0)}(x, y).$ 

Since  $f \perp g$ ,  $u \perp v$ ,  $a_0$  and  $w_1^{(0)}$  are *q*-proper polynomials, we may deduce that  $p(x, y) \mid g_1(x, qy)$ , i.e.,  $p(x, q^{-1}y) \mid g_1(x, y)$ . Let m(>0) be the greatest integer such that  $p(x, q^{-m}y) \mid g_1(x, y)$ . By virtue of (4.6), we have that  $p(x, q^{-m}y)$  divides

$$f(x, y)u_1(x, y)g_1(x, qy)v(x, qy)s(x, y)w_2^{(0)}(x, y)$$
  
 
$$\times \prod_{j=1}^{I} v(q^j x, y)w_2^{(j)}(x, y).$$

However,  $f \perp g$  and  $g_1 \perp u_1$  imply that  $p(x, q^{-m}y) \mid g_1(x, qy)$ , which contradicts the choice of *m*.

Case 2. Suppose  $p(x,qy) | g_1(x, y)$ . Let M > 0 be the greatest integer such that  $p(x,q^M y) | g_1(x, y)$ . Similarly, from (4.6) it follows that  $p(x,q^{M+1}y)$  divides

$$f(x,qy)r(x,y)u_1(x,qy)g_1(x,y)\prod_{j=0}^{I}v(q^jx,y)w_2^{(j)}(x,y)$$

Hence we get  $p(x, q^{M+1}y) | g_1(x, y)$ , which is again a contradiction.  $\Box$ 

To extend the above result to general bivariate *q*-hypergeometric terms, we need the concept of similar *q*-hypergeometric terms. Two bivariate *q*-hypergeometric terms  $T_1, T_2$  are called *similar* if there exists a rational function  $R \in \mathbb{F}(x, y)$  such that  $T_1(n, k)/T_2(n, k) = R(q^n, q^k)$ .

As in the ordinary case, the existence of qZ-pairs is preserved under the addition of similar bivariate q-hypergeometric terms.

**Lemma 4.5.** Suppose there exist qZ-pairs for two similar bivariate q-hypergeometric terms  $T_1(n, k)$  and  $T_2(n, k)$ . Then there exists a qZ-pair for  $T(n, k) = T_1(n, k) + T_2(n, k)$ .

Notice that T(n, k) = (K - 1)G(n, k) has a qZ-pair (1, G). Combining Theorem 4.4 and Lemma 4.5, we obtain the main result of this paper.

**Theorem 4.6.** Let T(n, k) be a bivariate q-hypergeometric term. Let  $T_1$ ,  $T_2$  be two similar bivariate q-hypergeometric terms satisfying

$$T(n,k) = (K-1)T_1(n,k) + T_2(n,k)$$

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and  $T_2(n, k)$  have a q-NR (r, s, u, v) with respect to K such that v is  $\varepsilon_y$ -free. Then T(n, k) has a qZ-pair if and only if  $T_2(n, k)$  is a q-proper hypergeometric term, or equivalently, if and only if v(x, y) is a q-proper polynomial.

#### 5. Algorithms

Let T(n, k) be a bivariate q-hypergeometric term. By the algorithm "q-RNF", we may find a q-NR (r, s, u, v) of T(n, k) with respect to K. Let

$$F(k) = \frac{u(x, q^k)}{v(x, q^k)} \prod_{j=0}^{k-1} \frac{r(x, q^j)}{s(x, q^j)}, \quad \forall k \in \mathbb{N}.$$

Then F(k) is a univariate q-hypergeometric term over the field  $\mathbb{F}(x)$  with a q-MR (r/s, u/v, 0). On the other hand, by Eq. (4.1), we have

$$\begin{aligned} \frac{F(k)|_{x=q^{n+1}}}{F(k)|_{x=q^n}} &= \frac{u(q^{n+1},q^k)v(q^n,q^k)}{u(q^n,q^k)v(q^{n+1},q^k)} \prod_{j=0}^{k-1} \frac{r(q^{n+1},q^j)s(q^n,q^j)}{r(q^n,q^j)s(q^{n+1},q^j)} \\ &= \frac{T(n+1,k)}{T(n,k)} \cdot \frac{T(n,0)}{T(n+1,0)} \cdot \frac{u(q^{n+1},q^0)v(q^n,q^0)}{u(q^n,q^0)v(q^{n+1},q^0)}, \end{aligned}$$

which is also a rational function of  $q^n$  and  $q^k$ . Hence  $\widetilde{F}(n,k) = F(k)|_{x=q^n}$  is a bivariate q-hypergeometric term.

Using the algorithm "q-decomp" given in Section 2, one may find univariate q-hypergeometric terms  $F_1(k)$ ,  $F_2(k)$  such that

$$F(k) = (K - 1)F_1(k) + F_2(k)$$

and  $F_2(k)$  has a q-MR  $(f_1/f_2, v_1/v_2, 0)$  with  $v_2$  being  $\boldsymbol{\varepsilon}_y$ -free. Since  $f_1/f_2, v_1/v_2 \in \mathbb{F}(x)(y)$ , we may assume that  $f_1, f_2, v_1, v_2 \in \mathbb{F}[x, y]$  and  $f_1 \perp f_2, v_1 \perp v_2$ . From the fact that r/s is  $\boldsymbol{\varepsilon}_y$ -reduced, it follows that  $f_1/f_2$  is also  $\boldsymbol{\varepsilon}_y$ -reduced.

Let

$$T_1(n,k) = T(n,0) \frac{v(q^n,q^0)}{u(q^n,q^0)} \cdot F_1(k)|_{x=q^n},$$
  
$$T_2(n,k) = T(n,0) \frac{v(q^n,q^0)}{u(q^n,q^0)} \cdot F_2(k)|_{x=q^n}.$$

Since Eq. (2.2) implies that

$$F_1(k) = \frac{U_1}{u/v} \cdot F(k)$$
 and  $F_2(k) = \frac{v_1/v_2}{u/v} \cdot F(k)$ ,

it follows that  $T_1(n, k)$  and  $T_2(n, k)$  are similar bivariate q-hypergeometric terms. It is easily verified that

$$T(n,k) = (K-1)T_1(n,k) + T_2(n,k)$$

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and  $(f_1, f_2, v_1, v_2)$  is a *q*-NR of  $T_2$  with respect to *K*. Therefore, Theorem 4.6 implies that T(n, k) has a *qZ*-pair if and only if  $v_2$  is a *q*-proper polynomial.

Finally, we need the algorithm given by Abramov and Le (2002) for determining whether or not a polynomial is q-proper.

We are now ready to describe the algorithm to determine whether a bivariate q-hypergeometric term T(n, k) has a qZ-pair.

1. Apply the algorithm in Böing and Koepf (1999) to find a rational function  $R \in \mathbb{F}(x, y)$  such that

$$\frac{T(n,k+1)}{T(n,k)} = R(q^n,q^k).$$

- 2. Find a q-RNF (r, s, u, v) with respect to  $\boldsymbol{\varepsilon}_{y}$  of R.
- 3. For D = r/s, U = u/v and  $n_0 = 0$ , apply the algorithm 'q-decomp' with respect to  $\boldsymbol{\varepsilon}_y$  to get  $V = v_1/v_2$ .
- 4. Use the algorithm in Abramov and Le (2002) to determine whether  $v_2$  is *q*-proper. If the answer is yes, then *T* has a *qZ*-pair; otherwise, *T* does not have any *qZ*-pair.

Here are two examples.

Example 1. Let

$$T(n,k) = \frac{q^k(1+q^{n+1}+q^{k+2})}{(q^n+q^k+1)(q^n+q^{k+1}+1)\prod_{j=1}^{k+1}(1-q^j)}$$

Then

$$\frac{T(n,k+1)}{T(n,k)} = \frac{q(1+q^{n+1}+q^{k+3})(q^n+q^k+1)}{(q^n+q^{k+2}+1)(1+q^{n+1}+q^{k+2})(1-q^{k+2})},$$

and we have

$$r = q, \ s = 1 - q^2 y, \ u = 1 + qx + q^2 y, \ v = (x + y + 1)(x + qy + 1)$$

is a q-NR of T with respect to K. For D = r/s, U = u/v and  $n_0 = 0$ , applying the algorithm "q-decomp", we get

$$V = v_1/v_2 = \frac{-q^2}{(-1+q^2)(x+1)}$$

Clearly,  $v_2$  is q-proper, so T(n, k) has a qZ-pair. Indeed, we can check that

$$L = 1, \quad G = \frac{1}{(q^n + q^k + 1) \prod_{j=1}^k (1 - q^j)}$$

is a qZ-pair for T(n, k).

Example 2.

$$T(n,k) = \frac{q^k(1+q^{n+1}+q^{k+2})}{(q^n+q^k+1)(q^n+q^{k+1}+1)\prod_{j=1}^k (1-q^j)}$$

Then

$$\frac{T(n,k+1)}{T(n,k)} = \frac{q(1+q^{n+1}+q^{k+3})(q^n+q^k+1)}{(q^n+q^{k+2}+1)(1+q^{n+1}+q^{k+2})(1-q^{k+1})},$$

and we have

$$r = q, \ s = 1 - qy, \ u = 1 + qx + q^2y, \ v = (x + y + 1)(x + qy + 1)$$

is a q-NR of T with respect to K. For D = r/s, U = u/v and  $n_0 = 0$ , applying the algorithm "q-decomp", we get

$$V = v_1/v_2 = \frac{-(x+y+1)q^2}{(q-1)(x+1)(x+qy+1)}$$

Since x + qy + 1 is not a q-proper polynomial, it follows that T(n, k) has no qZ-pair.

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