

CONGRUENCE CLASSES OF PRESENTATIONS FOR THE COMPLEX REFLECTION GROUPS $G(m, 1, n)$ AND $G(m, m, n)$

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ABSTRACT. In the present paper, we give a graph-theoretic description for representatives of all the congruence classes of presentations (or r.c.p. for brevity) for the imprimitive complex reflection groups $G(m, 1, n)$ and $G(m, m, n)$. We have three main results. The first main result is to establish a bijection between the set of all the congruence classes of presentations for the group $G(m, 1, n)$ and the set of isomorphism classes of all the rooted trees of n nodes. The next main result is to establish a bijection between the set of all the congruence classes of presentations for the group $G(m, m, n)$ and the set of isomorphism classes of all the connected graphs with n nodes and n edges. Then the last main result is to show that any generator set S of $G = G(m, 1, n)$ or $G(m, m, n)$ of n reflections, together with the respective basic relations on S , form a presentation of G .

In the paper [5], I introduced two concepts for any complex reflection group G generated by more than two reflections: one is the equivalence of simple root systems, and the other is the congruence of presentations. According to the definition, the equivalent simple root systems of G determine the congruent presentations of G . Then I, together with my students, Wang Li and Zeng Peng, found all the inequivalent simple root systems for all the primitive complex reflection groups except the group G_{34} (see [5][8][9]). We also described explicitly r.c.p. for the groups $G_{12}, G_{24}, G_{25}, G_{26}, G_7, G_{15}, G_{27}$. Then

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in [6], I further described explicitly r.c.p. for the groups G_{19}, G_{111}, G_{32} (here and later the notations for the complex reflection groups follow Shephard-Todd in [7], also see [1] [2]). In the present paper, we shall describe r.c.p. for the imprimitive complex reflection groups $G(m, 1, n)$ and $G(m, m, n)$. We first associate a reflection set X of $G(m, 1, n)$ to a graph Γ_X . Let $\Sigma(m, 1, n)$ (resp. $\Sigma(m, m, n)$) be the set of all the generator sets of $G(m, 1, n)$ (resp. $G(m, m, n)$) consisting of reflections with minimal possible cardinality. We show that any $X \in \Sigma(m, 1, n)$ consists of $n - 1$ reflections of type I and one reflection of order m with the graph Γ_X a tree of n nodes (see Lemma 2.1). Hence we can associate each $X \in \Sigma(m, 1, n)$ to a rooted tree of n nodes with the reflection of order m corresponding to the rooted node (see 1.6). On the other hand, we show that any $X \in \Sigma(m, m, n)$ consists of n reflections of type I with the number $\delta(X)$ (see 2.11 for the definition) coprime to m such that the graph Γ_X is connected and contains exactly one circle (see 2.10 and Theorem 2.19). Hence we can associate each $X \in \Sigma(m, m, n)$ to a connected graph with n nodes and exactly one circle. Then our first main result is to establish a bijection from the set of all the congruence classes of presentations for the group $G(m, 1, n)$ to the set of isomorphism classes of rooted trees of n nodes (see Theorem 3.2). Also, our second main result is to establish a bijection from the set of all the congruence classes of presentations of the group $G(m, m, n)$ to the set of isomorphism classes of connected graphs with n nodes and exactly one circle (see Theorem 3.4).

When a complex reflection group G is given, the congruence for a presentation (S, P) of G is entirely determined by the generator set S , not by the relation set P . However, in order to make a presentation of G more accessible, a proper choice of a relation set P is important in practice. To do this, we introduce the concept of basic relations on a generator set of the groups $G(m, 1, n)$ and $G(m, m, n)$. Then the third main result of the paper is to show that for $G = G(m, 1, n), G(m, m, n)$, any generator set S , together with the set P of basic relations on S form a presentation of G (see Theorems 4.18 and 4.21).

The technical tools in proving our main results are two operations on the reflection sets of $G(m, 1, n)$. One is called a terminal node operation, and the other is called a circle operation (see 2.5 and 2.11). Two facts are crucial in getting our results: one is that the set $\Sigma(m, 1, n)$ is transitive under the terminal node operations (2.5 and Lemma 2.6); the other is that the set $\Sigma(m, m, n)$ is transitive under the circle operations (2.11 and Lemma 2.14). These two facts enable us to reduce ourselves to certain simpler cases: for the group $G(m, 1, n)$, we reduce to the case where the graph Γ_X for some $X \in \Sigma(m, 1, n)$ is a string with the rooted node at one end; for $G(m, m, n)$, we reduce to the case where Γ_X for some $X \in \Sigma(m, m, n)$ is a string with a two-nodes circle at one end.

The contents are organized as follows. Section 1 is the preliminaries, some concepts and known results are collected there. In Sections 2 we describe the generator sets for the groups $G(m, 1, n)$ and $G(m, m, n)$. We establish a bijection between the congruence classes of presentations for each of these groups and the isomorphism classes of certain graphs in Section 3. Finally, we show that any generator set S of $G = G(m, 1, n)$ or $G(m, m, n)$, together with the respective basic relations on S , form a presentation of G .

§1. Preliminaries.

1.1. Let V be a complex vector space of dimension n . A *reflection* on V is a linear transformation on V of finite order with exactly $n-1$ eigenvalues equal to 1. A *reflection group* G on V is a finite group generated by reflections on V . A reflection group G on V is called a *real* group or a *Coxeter* group if there is a G -invariant \mathbb{R} -subspace V_0 of V such that the canonical map $\mathbb{C} \otimes_{\mathbb{R}} V_0 \rightarrow V$ is bijective. If this is not the case, G will be called *complex*. (Note that, according to this definition, a real reflection group is not complex.)

Since G is finite, there exists a unitary inner product $(\ , \)$ on V invariant under G . From now on we assume that such an inner product is fixed.

1.2. A reflection group G in V is *imprimitive*, if G acts on V irreducibly and there exists a decomposition $V = V_1 \oplus \dots \oplus V_r$ of nontrivial proper subspaces V_i , $1 \leq i \leq r$, of V such that G permutes $\{V_i \mid 1 \leq i \leq r\}$ (see [2]).

1.3. Let S_n be the symmetric group on n letters $1, 2, \dots, n$. For $\sigma \in S_n$, we denote by $[(a_1, \dots, a_n) | \sigma]$ the $n \times n$ monomial matrix with non-zero entries a_i in the $(i, (i)\sigma)$ -positions. For $p|m$ (read “ p divides m ”) in \mathbb{N} , we set

$$G(m, p, n) = \left\{ [(a_1, \dots, a_n) | \sigma] \mid a_i \in \mathbb{C}, a_i^m = 1; \left(\prod_j a_j \right)^{m/p} = 1; \sigma \in S_n \right\}$$

$G(m, p, n)$ is the matrix form of an imprimitive reflection group acting on V with respect to an orthonormal basis e_1, e_2, \dots, e_n , which is Coxeter only when either $m \leq 2$ or $(m, p, n) = (m, m, 2)$. We have $G(m, p, n) = G(1, 1, n) \times A(m, p, n)$, where $A(m, p, n)$ consists of all the diagonal matrices of $G(m, p, n)$, and $G(1, 1, n) \cong S_n$.

In particular, take $p = 1, m$, we get two special imprimitive reflection groups $G(m, 1, n)$ and $G(m, m, n)$ with $G(m, m, n) \subset G(m, 1, n)$. These two infinite families of groups are the main objects we shall study in the present paper.

In the present paper, we shall always assume $m, n > 2$ when we consider the groups $G(m, 1, n)$ or $G(m, m, n)$ unless otherwise specified.

1.4. For an orthonormal basis e_1, \dots, e_n of V and $\zeta_m = e^{2\pi i/m}$, let

$$\begin{aligned} R(m, m, n) &= \left\{ \frac{1}{\sqrt{2}} (\zeta_m^h e_i - \zeta_m^k e_j) \mid 1 \leq i \neq j \leq n, h, k \in \mathbb{Z} \right\} \\ R(m, n) &= \{ \zeta_m^k e_i \mid 1 \leq i \leq n, k \in \mathbb{Z} \}, \\ R(m, 1, n) &= R(m, m, n) \cup R(m, n). \end{aligned}$$

Then $R(m, m, n)$ and $R(m, 1, n)$ are *root systems* of the groups $G(m, m, n)$ and $G(m, 1, n)$ respectively, where the roots in $R(m, m, n)$ have order 2 and those in $R(m, n)$ have order m (see [2, 4.9] for the definition of a root system).

1.5. There are two kinds of reflections in the group $G(m, 1, n)$ as follows.

(i) One is with respect to a root in $R(m, m, n)$. It is of the form $s(i, j; k) = [(1, \dots, 1, \zeta_m^{-k}, 1, \dots, 1, \zeta_m^k, 1, \dots, 1)|(i, j)]$, where ζ_m^{-k}, ζ_m^k are the i th, resp. j th components of the n -tuple and (i, j) is the transposition of i and j for some $k \in \mathbb{Z}$ and $1 \leq i < j \leq n$. Call $s(i, j; k)$ a reflection of type I. Clearly, any reflection of type I has order 2. We also set $s(j, i; k) = s(i, j; -k)$ for any $1 \leq i < j \leq n$ and $k \in \mathbb{Z}$.

(ii) The other type of reflection is with respect to a root in $R(m, n)$. It is of the form $s(i; k) = [(1, \dots, 1, \zeta_m^k, 1, \dots, 1)|1]$ for some $k \in \mathbb{Z}$, where ζ_m^k occurs as the i th component of the n -tuple and 1 is the identity element of S_n . Call $s(i; k)$ a reflection of type II. $s(i; k)$ has order $m/\gcd(m, k)$.

All the reflections of type I lie in the subgroup $G(m, m, n)$.

1.6. For any $Z \subseteq \{1, 2, \dots, n\}$ with $r = |Z| > 0$, let V_Z be the subspace of V spanned by $\{e_i \mid i \in Z\}$. Let $R_Z(m, m, n) = R(m, m, n) \cap V_Z$ and $R_Z(m, 1, n) = R(m, 1, n) \cap V_Z$. Then $R_Z(m, m, n)$ (assuming $r > 1$) and $R_Z(m, 1, n)$ are root subsystems of $R(m, m, n)$ and $R(m, 1, n)$ respectively. Let $G_Z(m, 1, n)$ (resp. $G_Z(m, m, n)$) be the subgroup of $G(m, 1, n)$ (resp. $G(m, m, n)$) generated by the reflections with respect to the roots in $R_Z(m, 1, n)$ (resp. $R_Z(m, m, n)$). Then $G_Z(m, 1, n) \cong G(m, 1, r)$. When $r > 1$, we also have $G_Z(m, m, n) \cong G(m, m, r)$. To any set of reflections of $G_Z(m, 1, n)$ of type I, say $X = \{s(i_h, j_h; k_h) \mid h \in J\}$ for some index set J , we associate a digraph $\bar{\Gamma}_{Z, X} = (N_X, E_X)$ as follows. Its node set N_X is Z , and its arrow set E_X consists of all the ordered pairs (i, j) , $i < j$, with labels k for any $s(i, j; k) \in X$ (hence, if $s(i, j; k) \in X$ and $i > j$, then $\bar{\Gamma}_{Z, X} = (N_X, E_X)$ contains an arrow (j, i) with the label $-k$). Denote by $\Gamma_{Z, X}$ the underlying graph of $\bar{\Gamma}_{Z, X}$ which is obtained from $\bar{\Gamma}_{Z, X}$ by replacing all the labelled arrows by unlabelled edges.

Clearly, the graph $\Gamma_{Z, X}$ has no loop but may have multi-edges between two nodes.

The above definition of a graph can be extended: to any set X of reflections of $G_Z(m, 1, n)$, we define a graph $\Gamma_{Z, X}$ to be $\Gamma_{Z, X'}$, where X' is the subset of X consisting of all the reflections of type I. When X contains exactly one reflection of type II (say

$s(i; p)$), we define another graph, denoted by $\Gamma_{Z,X}^r$, which is obtained from $\Gamma_{Z,X}$ by rooting the node i , i.e., $\Gamma_{Z,X}^r$ is a rooted graph with the rooted node i . Sometimes we denote $\Gamma_{Z,X}^r$ by (Z, E_X, i) .

When $Z = \{1, 2, \dots, n\}$, we simply denote Γ_X (resp. Γ_X^r) for $\Gamma_{Z,X}$ (resp. $\Gamma_{Z,X}^r$).

Note that when S is the generator set in a presentation (G, S) of $G = G(m, 1, n), G(m, m, n)$, the graph defined here is different from a Coxeter-like graph given in [1, Appendix 2]: in a Coxeter-like graph, all the generating reflections are represented by nodes; while in a graph defined here, most of the generating reflections are represented by edges.

1.7. For a reflection group G , a *presentation of G by generators and relations* (or a *presentation* in short) is by definition a pair (S, P) , where

(1) S is a finite generator set for G which consists of reflections, and S has minimal cardinality with this property.

(2) P is a finite set of relations on S , and any other relation on S is a consequence of the relations in P .

A presentation (S, P) of G is *essential* if (S, P_0) is not a presentation of G for any proper subset P_0 of P .

Two presentations (S, P) and (S', P') for G are *congruent*, if there exists a bijection $\eta : S \rightarrow S'$ such that for any $s, t \in S$,

(*) $\langle s, t \rangle \cong \langle \eta(s), \eta(t) \rangle$, where the notation $\langle x, y \rangle$ stands for the subgroup generated by x, y .

In this case, we see by taking $s = t$ that the order $o(r)$ of r is equal to the order $o(\eta(r))$ of $\eta(r)$ for any $r \in S$.

If there does not exist such a bijection η , then we say that these two presentations are *non-congruent*.

When a reflection group G is a Coxeter group, the presentation of G as a Coxeter system is one presentation of G defined here. However, G may have some other presentations not congruent to the presentation of G as a Coxeter system. For example, let

G be the symmetric group S_n . Then one can show that the set of all the congruence classes of presentations of S_n is in one-to-one correspondence to the set of isomorphism classes of trees of n nodes. The presentation of S_n as a Coxeter system corresponds to the string of n nodes.

Given any reflection group G , by the above definition of a presentation, we see that for any generator set S of G with minimal possible cardinality, one can always find a relation set P on S such that (S, P) is a (essential) presentation of G . The congruence of the presentation (S, P) is entirely determined by the generator set S . So it makes sense to talk about the congruence relations among the generator sets of a reflection group G .

1.8. For any non-zero vector $v \in V$, denote by l_v the one dimensional subspace $\mathbb{C}v$ of V spanned by v , call it *a line*. In particular, denote $l_i = l_{e_i}$ for $1 \leq i \leq n$. Let $L = \{l_i \mid 1 \leq i \leq n\}$. Then the reflection $s(i, j; k)$ in $G(m, 1, n)$ interchanges the lines l_i, l_j and leaves all the other lines $l_h, h \neq i, j$, in L stable. The reflection $s(i; k)$ stabilizes all the lines in L . More generally, any element of $G(m, 1, n)$ gives rise to a permutation on the set L , and the action of $G(m, 1, n)$ (resp. $G(m, m, n)$) on L is transitive.

Let X be a set of reflections of $G(m, 1, n)$ and let $\langle X \rangle$ be the subgroup of $G(m, 1, n)$ generated by X . Then the action of $\langle X \rangle$ on L is transitive if and only if the graph Γ_X is connected. In particular, the graph Γ_X must be connected when X is the generator set of a presentation of $G(m, 1, n)$.

§2. The generator sets in the presentations for $G(m, 1, n)$ and $G(m, m, n)$.

In the present section, we shall describe the generator set S in a presentation (S, P) for the group $G(m, 1, n)$ or $G(m, m, n)$.

Lemma 2.1. *The generator set S of a presentation (S, P) of the group $G(m, 1, n)$ consists of $n - 1$ reflections of type I and one reflection of order m ($m, n > 2$ as assumed in 1.3). Hence the graph Γ_S is a tree.*

Proof. By the definition of a presentation and by [1, Appendix 2], the set S is of cardinality n . Let a_1 (resp. a_2) be the number of reflections in S of type I (resp. of order m). Then we have $a_2 \geq 1$ by the fact that any reflection of $G(m, 1, n)$ is conjugate to a power of some generating reflection (see [5, Section 1.5]). Hence $a_1 \leq n - 1$. So the graph Γ_S has at most $n - 1$ edges. Since the action of the group $G(m, 1, n)$ on L is transitive (see 1.8), the graph Γ_S must be connected and hence has at least $n - 1$ edges since it has n nodes. So $a_1 = n - 1$ and hence $a_2 = 1$. The last assertion follows immediately. \square

2.2. According to Lemma 2.1, we can define the graph Γ_S^r which is a rooted tree for any presentation (S, P) of the group $G(m, 1, n)$ (see 1.6). The generator set S consists of $n - 1$ reflections $s(i_h, j_h; k_h)$, $1 \leq h < n$, of type I and one reflection $s(p; k)$ for some $1 \leq p \leq n$ and $k \in \mathbb{Z}$ with k coprime to m (hence $o(s(p; k)) = m$, see 1.5 (ii)). By the fact that the graph Γ_S^r is a rooted tree, we have $\{i_h, j_h\} \neq \{i_l, j_l\}$ for any $h \neq l$. So for any $1 \leq h \neq l < n$, we have

$$\langle s(i_h, j_h; k_h), s(i_l, j_l; k_l) \rangle \cong \begin{cases} \mathbb{Z}_2 \times \mathbb{Z}_2, & \text{if } \{i_h, j_h\} \cap \{i_l, j_l\} = \emptyset, \\ S_3, & \text{if } \{i_h, j_h\} \cap \{i_l, j_l\} \neq \emptyset. \end{cases}$$

Also, we have

$$\langle s(i_h, j_h; k_h), s(p; k) \rangle \cong \begin{cases} \mathbb{Z}_2 \times \mathbb{Z}_m, & \text{if } p \notin \{i_h, j_h\}, \\ G(m, 1, 2), & \text{if } p \in \{i_h, j_h\}. \end{cases}$$

Here \mathbb{Z}_m is the cyclic group of order m .

2.3. Two rooted graphs (N, E, a) and (N', E', a') are *isomorphic*, if there exists a bijective map $\phi : N \rightarrow N'$ such that $\phi(a) = a'$ and that for any $v, w \in N$, $\{v, w\}$ is in E if and only if $\{\phi(v), \phi(w)\}$ is in E' .

The next result is concerned with a subgroup of $G(m, 1, n)$ generated by a set X of $n - 1$ reflections of type I with the graph Γ_X connected.

Lemma 2.4. *Let a_1, a_2, \dots, a_n be a permutation of $1, 2, \dots, n$ and let $X = \{s(a_h, a_{h+1}; k_h) \mid 1 \leq h < n\}$ be a reflection set of $G(m, 1, n)$ with some integers k_h . Then the subgroup $G = \langle X \rangle$ of $G(m, 1, n)$ generated by X is isomorphic to the symmetric group S_n .*

Proof. We know that $G(m, 1, n) = A(m, 1, n) \rtimes G(1, 1, n)$. The natural homomorphism $\pi : G(m, 1, n) \rightarrow G(1, 1, n) \cong S_n$ sends $s(a_h, a_{h+1}; k_h)$, $1 \leq h < n$, to $s_h = s(a_h, a_{h+1}; 0)$. Since s_1, s_2, \dots, s_{n-1} generate the group $G(1, 1, n)$, we have $\pi(G) \cong S_n$. Hence $|G| \geq |S_n|$, where $|X|$ denotes the cardinality of a set X . On the other hand, we have the following relations: for any $1 \leq h \leq l < n$,

$$(i) \ s(a_h, a_{h+1}; k_h)s(a_l, a_{l+1}; k_l) = s(a_l, a_{l+1}; k_l)s(a_h, a_{h+1}; k_h) \text{ if } l \neq h, h+1;$$

$$(ii) \ s(a_h, a_{h+1}; k_h)s(a_l, a_{l+1}; k_l)s(a_h, a_{h+1}; k_h) = s(a_l, a_{l+1}; k_l)s(a_h, a_{h+1}; k_h)s(a_l, a_{l+1}; k_l)$$

if $l = h+1$;

$$(iii) \ s(a_h, a_{h+1}; k_h)^2 = 1.$$

We know that the generator set $S = \{s_1, \dots, s_{n-1}\}$ together with the following set of relations:

$$(1) \ s_h s_l = s_l s_h \text{ if } l \neq h, h+1;$$

$$(2) \ s_h s_l s_h = s_l s_h s_l \text{ if } l = h+1;$$

$$(3) \ s_h^2 = 1.$$

form a presentation of S_n as a Coxeter system. This implies that G is also a homomorphic image of S_n , in particular, $|G| \leq |S_n|$. We get $|G| = |S_n|$. So the map $\pi : G \rightarrow S_n$ is an isomorphism. Our result follows. \square

2.5. Let X be a set of reflections of $G(m, 1, n)$ such that the graph Γ_X is connected and has two terminal nodes i, j (by a terminal node i of a graph, we mean that there is only one edge incident to i in the graph). Let $a_1 = i, a_2, \dots, a_r = j$ be a sequence of nodes such that X contains the following reflections: $t_h = s(a_h, a_{h+1}; k_h)$ for $1 \leq h < r$ and some integers k_h . Let $s = t_1 t_2 \dots t_{r-2} t_{r-1} t_{r-2} \dots t_1$. Then $s = s(a_1, a_r; k)$ with $k = \sum_{h=1}^{r-1} k_h$ is a reflection of $G(m, 1, n)$. Let $X' = (X \setminus \{t_{r-1}\}) \cup \{s\}$ (resp. $X'' = (X \setminus \{t_1\}) \cup \{s\}$). Then the graph $\Gamma_{X'}$ (resp. $\Gamma_{X''}$) is also a connected graph which can be obtained from Γ_X by replacing the edge $\{a_{r-1}, a_r\}$ (resp. $\{a_1, a_2\}$) by $\{a_1, a_r\}$. Call the transformation $\Gamma_X \mapsto \Gamma_{X'}$ (resp. $\Gamma_X \mapsto \Gamma_{X''}$) a *terminal node operation* on Γ_X . In this case, we see

that the graph Γ_X is a tree if and only if $\Gamma_{X'}$ (resp. $\Gamma_{X''}$) is a tree. Also, we have $\langle X' \rangle = \langle X \rangle$ (resp. $\langle X'' \rangle = \langle X \rangle$).

By abuse of terminology, we also say that the set X' (resp. X'') is obtained from X by a terminal node operation. In this case, it is not necessarily true that X is also obtained from X' (resp. X'') by a terminal node operation, it is the case only when the node a_{r-1} (resp. a_2) is a terminus in $\Gamma_{X'}$ (resp. $\Gamma_{X''}$).

A terminal node operation on X is applicable whenever the graph Γ_X is connected and contains at least two terminal nodes. In particular, this is the case when Γ_X is a tree ($n > 2$ as assumed).

Let X, Y be two sets of reflections in $G(m, 1, n)$. We say that X is obtained from Y (or equivalently, Y is transformed to X) by a sequence of terminal node operations, if there exists a sequence of reflection sets $X_1 = Y, X_2, \dots, X_r = X$ of $G(m, 1, n)$ with some $r \geq 1$ such that for every $1 < h \leq r$, X_h either is obtained from or gives rise to X_{h-1} by a terminal node operation.

Lemma 2.6. *Let X, Y be two subsets of $G(m, 1, n)$, each consisting of $n - 1$ reflections of type I such that the graphs Γ_X and Γ_Y are trees. Then Y can be transformed to some Y' by a sequence of terminal node operations such that the graphs Γ_X and $\Gamma_{Y'}$ are isomorphic.*

Proof. We may assume Γ_X is a string without loss of generality. If Γ_Y is also a string then we can just take $Y' = Y$. Now assume that we are not in that case. Fix a terminal node, say a_1 , of Γ_Y . Take another terminal node, say a_2 , of Γ_Y . Then we can apply a terminal node operation on Y with respect to a_1, a_2 to get Y_1 such that $\{a_1, a_2\}$ is an edge with a_2 a terminal node in Γ_{Y_1} . If Γ_{Y_1} is still not a string, we can take a terminal node $a_3 \neq a_2$ in Γ_{Y_1} . Applying a terminal node operation on Y_1 with respect to a_2, a_3 , we get Y_2 such that $\{a_2, a_3\}$ is an edge with a_3 a terminal node in Γ_{Y_2} . Continuing such a process, we can eventually get a reflection set Y_r for some $r > 1$ with Γ_{Y_r} a string by

a sequence of terminal node operations on Y . Then $Y' = Y_r$ is as required. \square

Note that for the reflection set X in Lemma 2.4, the graph Γ_X is a string. The following result generalizes Lemma 2.4.

Lemma 2.7. *Let X be a set of $n - 1$ reflections of $G(m, 1, n)$ such that the graph Γ_X is a tree. Then the subgroup $\langle X \rangle$ of $G(m, 1, n)$ generated by X is isomorphic to S_n .*

Proof. The tree Γ_X contains exactly $n - 1$ edge. Hence all the reflections in X have type I. If the graph Γ_X is a string then this is just the result of Lemma 2.4. Now assume that we are not in the case. Then by Lemma 2.6, we can transform X to some X' by a sequence of terminal node operations such that $\Gamma_{X'}$ is a string. Since $\langle X' \rangle = \langle X \rangle$, our result follows by Lemma 2.4. \square

The next result is the converse of Lemma 2.1.

Theorem 2.8. *Let X be a subset of $G(m, 1, n)$ consisting of $n - 1$ reflections of type I and one reflection of order m ($m > 2$ as assumed) such that the graph Γ_X is a tree. Then X generates the group $G(m, 1, n)$.*

Proof. Let X' be the set of $n - 1$ reflections of type I in X . Then the graph $\Gamma_{X'} = \Gamma_X$ is a tree. The natural map $\pi : G(m, 1, n) \rightarrow G(1, 1, n)$ sends $\langle X' \rangle$ isomorphically onto $G(1, 1, n)$ by Lemmas 2.4, 2.7 and their proof. The reflection of order m in X has the form $s(i; k)$ for some integers i, k with $1 \leq i \leq n$ and k coprime to m . Since the graph $\Gamma_{X'}$ is a tree and hence connected, the reflection $s(j; k)$ for any $1 \leq j \leq n$ can be obtained from $s(i; k)$ by $\langle X' \rangle$ -conjugation, hence $s(j; k) \in \langle X \rangle$. Now $A(m, 1, n)$ can be generated by all these $s(j; k)$'s and so it is contained in $\langle X \rangle$. This implies $\langle X \rangle \supseteq \langle X' \rangle A(m, 1, n) = G(m, 1, n) \supseteq \langle X \rangle$. Our result follows. \square

By Theorem 2.8, the following result can be used to show that the set of all the generator sets of $G(m, 1, n)$ are transitively permuted under the terminal node operations up to congruence.

Lemma 2.9. *Let X, Y be two subsets of $G(m, 1, n)$, each consisting of $n - 1$ reflections of type I and one reflection of order m such that the graphs Γ_X and Γ_Y are trees. So we can define rooted trees Γ_X^r and Γ_Y^r . Then Y can be transformed to some Y' by a sequence of terminal node operations such that the rooted trees Γ_X^r and $\Gamma_{Y'}^r$ are isomorphic.*

Proof. We may assume without loss of generality that Γ_X^r is a string with the rooted node at one end. By Lemma 2.6, we can transform Y to some Y'' by a sequence of terminal node operations such that $\Gamma_{Y''}$ is a string. Assume that a_1, a_2, \dots, a_n are nodes of $\Gamma_{Y''}^r$ with the node a_i rooted such that $\{a_h, a_{h+1}\}$, $1 \leq h < n$, are edges of $\Gamma_{Y''}^r$. If $i \in \{1, n\}$ then we can just take $Y' = Y''$. Otherwise, we can apply the terminal node operations on Y'' with respect to the node pairs $\{a_1, a_n\}, \{a_2, a_1\}, \dots, \{a_{i-1}, a_{i-2}\}$ in turn. Then the result is just a required set Y' . \square

2.10. By Lemma 2.7 and [1, Appendix 2], we see that any presentation (S, P) of $G(m, m, n)$ must satisfy $|S| = n$. That is, Γ_S is a connected graph with n nodes and n edges. So Γ_S must be a connected graph with exactly one circle, where we allow a circle to have just two nodes, i.e., a pair of nodes with double edges. In the remaining part of this section, we shall give a necessary and sufficient condition for a set X of n reflections with Γ_X connected to generate the group $G(m, m, n)$. To do this, we shall first introduce a new operation on X .

2.11. Assume that X is a reflection set of $G(m, m, n)$ such that Γ_X is connected and contains exactly one circle, say the edges of the circle are $\{a_h, a_{h+1}\}$, $1 \leq h \leq r$ (the subscripts are modulo r) for some integer $2 \leq r \leq n$. Then X contains the reflections $s(a_h, a_{h+1}; k_h)$ with some integers k_h for any $1 \leq h \leq r$ (the subscripts are modulo r). Denote by $\delta(X)$ the absolute value of $\sum_{h=1}^r k_h$.

Suppose that Γ_X also contains an edge $\{a_0, a_1\}$ with $a_0 \neq a_2, a_r$. Hence X contains a reflection $s(a_0, a_1; k_0)$ for some $k_0 \in \mathbb{Z}$. Let $Y = (X \setminus \{s(a_r, a_1; k_r)\}) \cup \{s(a_r, a_0; k_r - k_0)\}$. Then the graph Γ_Y can be obtained from Γ_X by replacing the edge $\{a_r, a_1\}$ by

$\{a_r, a_0\}$. Clearly, the graph Γ_Y is also connected and contains exactly one circle with $\delta(Y) = \delta(X)$. Call the transformation $X \mapsto Y$ a *circle expansion* and call the reverse transformation $Y \mapsto X$ a *circle contraction*. Call both transformations *circle operations*. Since $s(a_r, a_0; k_r - k_0) = s(a_0, a_1; k_0)s(a_r, a_1; k_r)s(a_0, a_1; k_0)$, we have $\langle Y \rangle = \langle X \rangle$.

Clearly, a circle contraction on X is applicable whenever X has a circle with at least three nodes. Also, a circle expansion on X is applicable whenever Γ_X contains a circle which is incident to some edge at one node.

2.12. Let Δ be the set of all the reflection sets X of $G(m, m, n)$ with Γ_X connected and containing exactly one circle. Then all applicable circle operations stabilize the set Δ . The fibres of the function $\delta : \Delta \rightarrow \mathbb{Z}_{\geq 0}$ (the set of non-negative integers) are stable under the circle operations. Assume that Y is obtained from $X \in \Delta$ by a sequence of circle operations. Then $\langle Y \rangle = \langle X \rangle$. Hence we get

Lemma 2.13. *Let X and Y be in Δ such that Y can be obtained from X by a sequence of circle operations. Then we have $\langle Y \rangle = \langle X \rangle$ and $\delta(Y) = \delta(X)$.*

Lemma 2.14. *Let X be a reflection set of $G(m, m, n)$ such that the graph Γ_X is connected and contains exactly one circle. Then X can be transformed to some X' by a sequence of circle operations such that the graph $\Gamma_{X'}$ is a string with a two-nodes circle at one end.*

Proof. First we can apply a sequence of circle expansions to transform X to some X'' such that the graph $\Gamma_{X''}$ is a circle. Let c_1, c_2, \dots, c_n be the nodes of $\Gamma_{X''}$ such that X'' consists of the reflections $t_h = s(c_h, c_{h+1}; k_h)$ for $1 \leq h \leq n$ (the subscripts are modulo n). Let $t'_{n-1} = t_n t_{n-1} t_n$ and $t'_j = t'_{j+1} t_j t'_{j+1}$ for $2 \leq j < n-1$. Let $X_{n-1} = (X'' \setminus \{t_n\}) \cup \{t'_{n-1}\}$ and $X_j = (X_{j+1} \setminus \{t'_{j+1}\}) \cup \{t'_j\}$ for $2 \leq j < n-1$. Denote $X_n = X''$. Then X_j is obtained from X_{j+1} by a circle contraction and $X_j \in \Sigma(m, m, n)$ for $2 \leq j < n$. Clearly, $X' = X_2$ is a required element in $\Sigma(m, m, n)$. \square

2.15. Let $X = \{s(h, h+1; k_h), s(1, 2; k'_1) \mid 1 \leq h < n\}$ for some integers $k'_1, k_1, \dots, k_{n-1}$.

Then the graph Γ_X is a string with a two-nodes circle at one end. We want to describe the subgroup $\langle X \rangle$ of $G(m, m, n)$ generated by X .

Denote $t_h = s(h, h+1; k_h)$, $1 \leq h < n$, and $t'_1 = s(1, 2; k'_1)$. Let $\zeta = e^{2\pi i/m}$. For any $1 < i \leq n$ and $k \in \mathbb{Z}$, denote $\alpha(i; k) = [(\zeta^{-k}, 1, \dots, 1, \zeta^k, 1, \dots, 1)|1]$, where ζ^k is the i th component of the n -tuple. Then we have $t'_1 t_1 = \alpha(2; k'_1 - k_1)$. Denote $p = k'_1 - k_1$. We have $\alpha(i; ap) = \alpha(i; p)^a = t_{i-1} t_{i-2} \dots t_2 (t'_1 t_1)^a t_2 \dots t_{i-1}$ for any $2 \leq i \leq n$ and $a \in \mathbb{Z}$. For $c_2, \dots, c_n \in \mathbb{Z}$, denote $\beta(c_2, \dots, c_n) = [(\zeta^{c_2 p}, \zeta^{c_2 p}, \dots, \zeta^{c_n p})|1]$, where $c = -\sum_{i=2}^n c_i$. Then $\beta(c_2, \dots, c_n) = \prod_{i=2}^n \alpha(i; c_i p) = \prod_{i=2}^n \alpha(i; p)^{c_i}$. Let $N = \{\beta(c_2, \dots, c_n) \mid c_i \in \mathbb{Z}\}$. Clearly, N is closed under multiplication and inversion. So N itself forms an abelian group. By the above discussion, we see that N is a subgroup of $\langle X \rangle$. For any $c_2, \dots, c_n \in \mathbb{Z}$ and $t \in X$, we have $\beta(c'_2, \dots, c'_n) = t\beta(c_2, \dots, c_n)t$, where, if $t = t_h$ with $2 \leq h < n$ then the sequence c'_2, \dots, c'_n can be obtained from c_2, \dots, c_n by transposing the terms c_h and c_{h+1} ; if $t \in \{t'_1, t_1\}$ then

$$c'_i = \begin{cases} c_i, & \text{if } i > 2, \\ -\sum_{j=2}^n c_j, & \text{if } i = 2. \end{cases}$$

This implies that N is an abelian normal subgroup of $\langle X \rangle$.

Lemma 2.16. *In the above setup, we have $\langle X \rangle = N \rtimes \langle t_1, \dots, t_{n-1} \rangle$.*

Proof. We have just shown that $N \triangleleft \langle X \rangle$. We know by Lemma 2.4 and its proof that $\langle t_1, \dots, t_{n-1} \rangle$ is a subgroup of $\langle X \rangle$ and that the natural map $G(m, m, n) \rightarrow G(1, 1, n)$ sends $\langle t_1, \dots, t_{n-1} \rangle$ isomorphically onto $G(1, 1, n)$. So the intersection of N and $\langle t_1, \dots, t_{n-1} \rangle$ is the trivial subgroup $\{1\}$. Hence $N\langle t_1, \dots, t_{n-1} \rangle = N \rtimes \langle t_1, \dots, t_{n-1} \rangle$, which is a subgroup of $\langle X \rangle$. It remains to show that $\langle X \rangle \subseteq N \rtimes \langle t_1, \dots, t_{n-1} \rangle$. To do this, it is enough to show that $t \cdot (N \rtimes \langle t_1, \dots, t_{n-1} \rangle) \subseteq N \rtimes \langle t_1, \dots, t_{n-1} \rangle$ for any $t \in X$. It is also enough to show that $t\beta(c_2, \dots, c_n) \subseteq N \rtimes \langle t_1, \dots, t_{n-1} \rangle$ for any $t \in X$ and $c_i \in \mathbb{Z}$. When $t = t_h$ for some $1 \leq h < n$, we have $t\beta(c_2, \dots, c_n) = (t_h\beta(c_2, \dots, c_n)t_h)t_h \subseteq Nt_h$; when $t = t'_1$, we have $t\beta(c_2, \dots, c_n) = \alpha(2; p)t_1\beta(c_2, \dots, c_n) = (\alpha(2; p)t_1\beta(c_2, \dots, c_n)t_1)t_1 \subseteq Nt_1$ since $\alpha(2; p) = \beta(1, 0, \dots, 0) \in N$. This proves our result. \square

Remark 2.17. The above lemma can be generalized: The set X could be taken a more general reflection set $X' = \{s(a_h, a_{h+1}; k_h), s(a_1, a_2; k'_1) \mid 1 \leq h < n\}$, where $k'_1, k_1, \dots, k_{n-1}$ are any fixed integers, and a_1, \dots, a_n is any permutation of $1, 2, \dots, n$. Actually, such a more general subgroup $\langle X' \rangle$ of $G(m, m, n)$ is T -conjugate to a subgroup of the form $\langle X \rangle$ in Lemma 2.16, where T is the permutation matrix in $G(m, m, n)$ corresponding to the permutation a_1, a_2, \dots, a_n . Hence we also have a decomposition $\langle X' \rangle = N' \rtimes \langle s_1, \dots, s_{n-1} \rangle$, where $s_h = s(a_h, a_{h+1}; k_h)$ for $1 \leq h < n$, and $N' = \langle X' \rangle \cap A(m, m, n) = T \langle X \rangle T^{-1} \cap A(m, m, n) = T N T^{-1} = N$.

Corollary 2.18. *Let $X' = \{s(a_h, a_{h+1}; k_h), s(a_1, a_2; k'_1) \mid 1 \leq h < n\}$, where $k'_1, k_1, \dots, k_{n-1}$ are any fixed integers, and a_1, \dots, a_n is any permutation of $1, 2, \dots, n$. Then X' is a generator set of $G(m, m, n)$ if and only if the integer $k'_1 - k_1$ is coprime to m .*

Proof. By the above remark, we need only consider the case when $a_h = h$ for any $1 \leq h \leq n$. Hence X' is just the set X in Lemma 2.16. By Lemma 2.16, we see that $\langle X \rangle = G(m, m, n)$ if and only if the subgroup N of $\langle X \rangle$ in Lemma 2.16 is equal to $A(m, m, n)$. The latter holds if and only if the number $p = k'_1 - k_1$ in 2.15 is coprime to m . Hence the result. \square

Note that in the setup of Corollary 2.18, the number $|k'_1 - k_1|$ is equal to $\delta(X')$.

Now we are ready to give a criterion for a certain reflection set to be a generator set of $G(m, m, n)$.

Theorem 2.19. *Let X be a reflection set of $G(m, m, n)$ such that the graph Γ_X is connected and contains exactly one circle. Then X generates $G(m, m, n)$ if and only if the integer $\delta(X)$ is coprime to m .*

Proof. By Lemma 2.14, we can transform X to some reflection set X' by a sequence of circle operations such that the graph $\Gamma_{X'}$ is a string with a two-nodes circle at one end. Then by Lemma 2.13, we have $\langle X' \rangle = \langle X \rangle$ and $\delta(X') = \delta(X)$. Hence by Corollary

2.18, we see that $\langle X' \rangle = G(m, m, n)$ if and only if $\delta(X')$ is coprime to m . So our result follows immediately. \square

§3. The congruence classes of presentations of the group $G(m, 1, n)$.

In the present section, we shall get two main results of the paper, which establish a bijection from the set of all the congruence classes of presentations for the group $G(m, 1, n)$ (resp., $G(m, m, n)$) to the set of isomorphism classes of certain connected graphs.

3.1. Let S be any generator set of $G(m, 1, n)$ consisting of n reflections. By Lemma 2.1 and Theorem 2.8, we see that S gives rise to a presentation (S, P) of the group $G(m, 1, n)$ for a certain set P of relations on S . By Lemma 2.1, $S = S' \cup \{s\}$, where S' consists of $n - 1$ reflections of type I with the graph $\Gamma_{S'}$ being a tree, and s is a reflection of order m . So we can define the rooted tree Γ_S^r . The set S satisfies the following relations:

- (1) $t^2 = 1$ for any $t \in S'$;
- (2) $s^m = 1$;
- (3) $tt' = t't$ for any $t, t' \in S'$ with the edges $e(t), e(t')$ having no common node in Γ_S^r ;
- (4) $tt't = t'tt'$ for any $t, t' \in S'$ with $e(t), e(t')$ having exactly one common node in Γ_S^r ;
- (5) $ts = st$ for any $t \in S'$ with the edge $e(t)$ not incident to the rooted node $n(s)$;
- (6) $tsts = stst$ for any $t \in S'$ with the edge $e(t)$ incident to the rooted node $n(s)$.

Call all the above relations in (1)-(6) *the order and braid relations (o.b. relations in short)* on S . Suppose that P is a certain relation set on S such that (S, P) forms a presentation of $G(m, 1, n)$. Then the congruence class of (S, P) is entirely determined by the generator set S and the o.b. relations on S , the latter is entirely determined by the isomorphism class of the rooted tree Γ_S^r up to congruence. In other words, two

presentations (S_1, P_1) and (S_2, P_2) of $G(m, 1, n)$ are congruent if and only if the rooted trees $\Gamma_{S_1}^r$ and $\Gamma_{S_2}^r$ are isomorphic. On the other hand, let $\Gamma^r = ([n], E, i)$ be a rooted tree, where $i \in [n] = \{1, 2, \dots, n\}$. Then $|E| = n - 1$. We can define a reflection set X of $G(m, 1, n)$ as follows. The reflection $s(i, j; 0)$ is assigned to X if and only if $\{i, j\} \in E$. Also, the reflection $s(i; 1)$ is assigned to X . Hence X consists of $n-1$ reflections $s(i, j; 0)$, $\{i, j\} \in E$, of type I and one reflection $s(i; 1)$ of order m . By Theorem 2.8, we see that X forms a generator set of $G(m, 1, n)$ with $\Gamma_X^r = \Gamma^r$. So we get the following

Theorem 3.2. *The map $(S, P) \rightarrow \Gamma_S^r$ induces a bijection from the set of all the congruence classes of presentations for the group $G(m, 1, n)$ to the set of isomorphism classes of rooted trees with n nodes.*

3.3. Next let S be any generator set of $G(m, m, n)$ consisting of n reflections. By 2.10 and Lemma 2.7, we see that S gives rise to a presentation (S, P) of the group $G(m, m, n)$ for a certain set P of relations on S . By 2.10 and Theorem 2.19 we see that the graph Γ_S is connected and contains exactly one circle with $\delta(S)$ coprime to m . The set S satisfies the following relations:

- (1) $t^2 = 1$ for any $t \in S$;
- (2) $st = ts$ for any $s, t \in S$ with the edges $e(s), e(t)$ having no common node in Γ_S ;
- (3) $sts = tst$ for any $s, t \in S$ with $e(s), e(t)$ having exactly one common node in Γ_S ;

When Γ_S has a two-nodes circle, say $s, t \in S$ with $e(s), e(t)$ two edges of the circle, we have

- (4) $(st)^m = 1$.

Call all the above relations in (1)-(4) *the order and braid relations (o.b. relations in short)* on S . Suppose that P is a certain relation set on S such that (S, P) forms a presentation of $G(m, m, n)$. Then the congruence class of (S, P) is entirely determined by the generator set S and the o.b. relations on S , the latter is entirely determined by the isomorphism class of the graph Γ_S up to congruence. In other words, two presentations

(S_1, P_1) and (S_2, P_2) of $G(m, m, n)$ are congruent if and only if the graphs Γ_{S_1} and Γ_{S_2} are isomorphic. On the other hand, let $\Gamma = ([n], E)$ be a connected graph with n nodes and exactly one circle. Then $|E| = n$. Fix any $\{p, q\} \in E$ in the circle of Γ . We can define a reflection set X of $G(m, m, n)$ as follows. The reflection $s(i, j; 0)$ is assigned to X for all $\{i, j\} \in E \setminus \{\{p, q\}\}$. Also, the reflection $s(p, q; 1)$ is assigned to X . Hence X consists of n reflections with $\delta(X) = 1$, coprime to m . By Theorem 2.19, we see that X forms a generator set of $G(m, m, n)$ with $\Gamma_X = \Gamma$. So we get the following

Theorem 3.4. *The map $(S, P) \rightarrow \Gamma_S$ induces a bijection from the set of all the congruence classes of presentations for the group $G(m, m, n)$ to the set of isomorphism classes of connected graphs with n nodes and n edges (or equivalently with n nodes and exactly one circle).*

§4. The relation sets of the presentations for the groups $G(m, 1, n)$ and $G(m, m, n)$.

For any generator set S of the group $G(m, 1, n)$ (resp. $G(m, m, n)$) of n reflections, we can always find a set P of relations on S such that (S, P) is a presentation of $G(m, 1, n)$ (resp. $G(m, m, n)$). However, a relation set P is not uniquely determined by S . Then the problem is how to choose a relation set P such that (S, P) becomes a presentation of $G(m, 1, n)$ (resp. $G(m, m, n)$). We shall find one for each generator set of the group.

4.1. It is well known that the group $G(m, 1, n)$ has a presentation (S, P) , where $S = \{s(h, h+1; 0), s(1; 1) \mid 1 \leq h < n\}$, and P consists of the following relations: denote $t_h = s(h, h+1; 0)$, $1 \leq h < n$, and $t_0 = s(1; 1)$,

- (i) $t_0^m = 1$;
- (ii) $t_h^2 = 1$ for $1 \leq h < n$;
- (iii) $t_i t_j = t_j t_i$ if $j \neq i \pm 1$;
- (iv) $t_i t_{i+1} t_i = t_{i+1} t_i t_{i+1}$ for $1 \leq i < n-1$;
- (v) $t_0 t_i = t_i t_0$ for $i > 1$;

$$(vi) \ t_0 t_1 t_0 t_1 = t_1 t_0 t_1 t_0.$$

4.2. Let $\Sigma(m, 1, n)$ be the set of all the generator sets S of $G(m, 1, n)$ of n reflections such that the graph Γ_S is a tree (hence S consists of $n - 1$ reflections of type I and one reflection of order m by Lemma 2.1). We can define a rooted tree Γ_S^r for any $S \in \Sigma(m, 1, n)$ by 1.6. We know that the set $\Sigma(m, 1, n)$ is stable under terminal node operations. By Lemma 2.9, we see that for any X, X' in $\Sigma(m, 1, n)$, one can find a sequence $X_1 = X', X_2, \dots, X_r = X''$ in $\Sigma(m, 1, n)$ such that for every $1 < h \leq r$, X_h either is obtained from or gives rise to X_{h-1} by a terminal node operation and that X is congruent to X'' (see 1.7). In this sense, one can say that the terminal node operations act on the congruence classes of $\Sigma(m, 1, n)$ transitively.

4.3. Let $X = \{s(i_h, j_h; k_h), s(l; k) \mid 1 \leq h < n\}$ be in $\Sigma(m, 1, n)$. Then k_h, k, l, i_h, j_h are some integers with $1 \leq l \leq n$, $1 \leq i_h \neq j_h \leq n$, and k coprime to m . The following relations on X hold: denote $s_h = s(i_h, j_h; k_h)$, $1 \leq h < n$, and $s = s(l; k)$,

$$(1) \ s^m = 1;$$

$$(2) \ s_h^2 = 1 \text{ for } 1 \leq h < n;$$

$$(3) \ s_p s_q = s_q s_p \text{ if the edges } e(s_p) \text{ and } e(s_q) \text{ have no common node};$$

$$(4) \ s_p s_q s_p = s_q s_p s_q \text{ if } e(s_p) \text{ and } e(s_q) \text{ have exactly one common node};$$

$$(5) \ s s_p = s_p s \text{ if the edge } e(s_p) \text{ is not incident to the node } l;$$

$$(6) \ s s_p s s_p = s_p s s_p s \text{ if } e(s_p) \text{ is incident to } l;$$

$$(7) \ s \cdot s_p s_q s_p = s_p s_q s_p \cdot s \text{ if } e(s_p) \text{ and } e(s_q) \text{ have exactly one common node } l;$$

$$(8) \ s_p \cdot s_q s_r s_q = s_q s_r s_q \cdot s_p \text{ if } s_p, s_q, s_r \text{ are pairwise distinct with the edges } e(s_p), e(s_q)$$

and $e(s_r)$ incident to one common node.

Call the relations (1)-(2) *order relations*, (3)-(6) *braid relations*, (7) *root-braid relations*, and (8) *branching relations* on X . Call all of them *basic relations* on X .

4.4. Let $\Sigma(m, m, n)$ be the set of all the generator sets S of the group $G(m, m, n)$ each of which consists of n reflections. Given any $S \in \Sigma(m, m, n)$, we see by 2.10 and Theorem 2.19 that the graph Γ_S is connected and contains exactly one circle. Hence $\Sigma(m, m, n)$

is a subset of the set Δ defined in 2.12. By Lemma 2.13 and Theorem 2.19, we see that $\Sigma(m, m, n)$ is stable under the circle operations. For any $S \in \Sigma(m, m, n)$, we can find a relation set P on S such that (S, P) is a presentation of $G(m, m, n)$. The congruence class of (S, P) is entirely determined by the generator set S and the o.b. relations on the set S , the latter is determined by the isomorphism class of the graph Γ_S up to congruence. So we can talk about the congruence of any $S \in \Sigma(m, m, n)$. By Lemma 2.14, we can find, for any two $S, S' \in \Sigma(m, m, n)$, a sequence $S_1 = S', S_2, \dots, S_r = S''$ in $\Sigma(m, m, n)$ with S, S'' congruent such that for every $1 < h \leq r$, S_h is obtained from S_{h-1} by a circle operation. In this sense, we can say that the set $\Sigma(m, m, n)$ is transitive under circle operations.

4.5. For $S \in \Sigma(m, m, n)$, let c_1, c_2, \dots, c_r be the nodes of the circle in Γ_S such that $\{c_i, c_{i+1}\}$, $1 \leq i \leq r$, are edges of Γ_S , where the subscripts are modulo r . Hence S contains the reflections $t_i = s(c_i, c_{i+1}; p_i)$ for $1 \leq i \leq r$ (the subscripts are modulo r) and some integers p_i . Let $t_j = s(a_j, b_j; k_j)$, $r < j \leq n$, be the reflections such that $S = \{t_i \mid 1 \leq i \leq n\}$. For $1 \leq i < j \leq r$, denote $s_{ij} = t_i \dots t_{j-2} t_{j-1} t_{j-2} \dots t_i$ and $s_{ji} = t_{i-1} \dots t_1 t_r \dots t_{j+1} t_j t_{j+1} \dots t_r t_1 \dots t_{i-1}$. In particular, $s_{ij} = t_i$ if $j = i + 1$ and $s_{ji} = t_j$ if $j = r$ and $i = 1$. Then the following relations on S hold:

- (i) $t_i^2 = 1$ for $1 \leq i \leq n$;
- (ii) $t_i t_j = t_j t_i$ if the edges $e(t_i), e(t_j)$ corresponding to t_i, t_j have no common node;
- (iii) $t_i t_j t_i = t_j t_i t_j$ if the edges $e(t_i), e(t_j)$ have exactly one common node;
- (iv) $(s_{ij} s_{ji})^m = 1$ for any $1 \leq i < j \leq r$;
- (v) $s \cdot trt = trt \cdot s$ for any branching node b of Γ_S and any triple $X = \{s, t, r\} \subseteq S$

with Γ_X having b as its branching node;

- (vi) (a) $us_{ij}u \cdot vs_{ji}v = vs_{ji}v \cdot us_{ij}u$,
- (b) $us_{ij}s_{ji}us_{ij}s_{ji} = s_{ij}s_{ji}us_{ij}s_{ji}u$, and
- (c) $vs_{ij}s_{ji}vs_{ij}s_{ji} = s_{ij}s_{ji}vs_{ij}s_{ji}v$

if there are some $u, v \in S$ with $e(u), e(v)$ incident to the circle of Γ_S at the nodes c_i, c_j

respectively for some $1 \leq i < j \leq r$.

Call relations (i) *order relations*, call (ii)-(iii) *braid relations*, call (iv) *circle relations*. (iv) in case of $r = 2$ is also called *a braid relation*, Call (v) *branching relations*, and call (vi) *branching-circle relations*. Call all of these relations *basic relations* on S .

Remark 4.6. (1) When $j - i \geq 2$, we have $t_i(s_{ij}s_{ji})t_i = s_{i+1,j}s_{j,i+1}$ and $t_j(s_{ij}s_{ji})t_j = s_{i,j+1}s_{j+1,i}$. Thus relation $(s_{ij}s_{ji})^m = 1$ holds if and only if $(s_{i+1,j}s_{j,i+1})^m = 1$ holds if and only if $(s_{i,j+1}s_{j+1,i})^m = 1$ holds. So in (iv), we need only write down one of such relations for a fixed pair $1 \leq i < j \leq r$.

(2) We shall show in Lemmas 4.8 and 4.10 that the branching and root-braid relations can be deduced from a certain part of these relations.

(3) Note that relation (b) (resp. (c)) in (vi) holds whenever such a reflection u (resp. v) exists. For some smaller number r , a computer programme, called *MAGMA*, shows that when both u and v exist, relation (a) in (vi) implies relations (b) and (c). It is natural to conjecture that this should hold for any positive integer r .

4.7. Let $S \in \Sigma(m, 1, n)$ (resp. $S \in \Sigma(m, m, n)$) be with c a branching node in the graph Γ_S . Let c_1, \dots, c_r be all distinct nodes with $r > 2$ such that $\{c_i, c\}$, $1 \leq i \leq r$, are edges of Γ_S . Then S contains reflections $t_i = s(c_i, c; k_i)$ for some integers k_i . We have the following relations:

- (i) $t_i t_j t_i = t_j t_i t_j$ and $t_i^2 = 1$ for any $1 \leq i \neq j \leq r$;
- (ii) $t_l \cdot t_i t_j t_i = t_i t_j t_i \cdot t_l$ for any distinct i, j, l in $\{1, 2, \dots, r\}$.

The following result shows that the branching relations (7) on $S \in \Sigma(m, 1, n)$ (resp. the branching relations (v) on $S \in \Sigma(m, m, n)$) can be deduced from some part of these relations.

Lemma 4.8. *Fix p , $1 \leq p \leq r$. Under the assumption of condition (i), condition (ii) is equivalent to the following condition*

$$(i') \quad s_p \cdot s_i s_j s_i = s_i s_j s_i \cdot s_p \text{ for any } 1 \leq i \neq j \leq r \text{ with } i, j \neq p.$$

Proof. It is clear that (ii) implies (ii'). Now assume (ii'). We want to show (ii). It is easily seen by condition (i) that relation $s_l \cdot s_i s_j s_i = s_i s_j s_i \cdot s_l$ holds in the case when $p \in \{i, j, l\}$. It remains to show the relation $s_l \cdot s_i s_j s_i = s_i s_j s_i \cdot s_l$ in the case of $p \notin \{i, j, l\}$. We have $s_l \cdot s_p s_i s_p = s_p s_i s_p \cdot s_l$, $s_l \cdot s_p s_j s_p = s_p s_j s_p \cdot s_l$ and $s_p \cdot s_i s_j s_i = s_i s_j s_i \cdot s_p$. Hence $s_l \cdot s_p s_i s_j s_i s_p = s_l \cdot s_p s_i s_p \cdot s_p s_j s_p \cdot s_p s_i s_p = s_p s_i s_p \cdot s_p s_j s_p \cdot s_p s_i s_p \cdot s_l = s_p s_i s_j s_i s_p \cdot s_l$. Then $s_p s_l s_p \cdot s_i s_j s_i = s_i s_j s_i \cdot s_p s_l s_p$. Hence $s_p s_l \cdot s_i s_j s_i s_p = s_p s_i s_j s_i \cdot s_l s_p$. This implies $s_l \cdot s_i s_j s_i = s_i s_j s_i \cdot s_l$. \square

4.9. Under the setup of 4.7 with $S \in \Sigma(m, 1, n)$, assume $s = s(c; k) \in S$. Then the root-braid relations on S are

$$(iii) \quad s \cdot t_i t_j t_i = t_i t_j t_i \cdot s \text{ for any distinct } i, j \text{ in } \{1, 2, \dots, r\}.$$

The following result shows that the root-braid relations on $S \in \Sigma(m, 1, n)$ can be deduced from some part of the relations.

Lemma 4.10. *Fix p , $1 \leq p \leq r$. Under the assumption of conditions 4.7 (i)-(ii), condition (iii) is equivalent to the following condition*

$$(iii') \quad s \cdot s_p s_j s_p = s_p s_j s_p \cdot s \text{ for any } 1 \leq j \leq r \text{ with } j \neq p.$$

Proof. It is clear that (iii) implies (iii'). Now assume (iii'). We want to show (iii). We must show $s \cdot s_i s_j s_i = s_i s_j s_i \cdot s$ for any $1 \leq i \neq j \leq r$ with $i, j \neq p$. We have $s \cdot s_p s_i s_p \cdot s_p s_j s_p \cdot s_p s_i s_p = s_p s_i s_p \cdot s_p s_j s_p \cdot s_p s_i s_p \cdot s$ by (iii'). Hence $s \cdot s_p s_i s_j s_i s_p = s_p s_i s_j s_i s_p \cdot s$ by 4.7 (i). This implies $s \cdot s_i s_j s_i = s_i s_j s_i \cdot s$ by 4.7 (ii). \square

4.11. Suppose that $X \in \Sigma(m, m, n)$ contains the reflections $t_h = s(c_h, c_{h+1}; k_h)$ (the subscripts are modulo r) for $1 \leq h \leq r$ and some nodes c_1, \dots, c_r of Γ_X with $r > 2$. Suppose that for some $1 \leq i < j \leq r$, there are some $s, t \in X \setminus \{t_{i-1}, t_i, t_{j-1}, t_j\}$ with the edge $e(s)$ incident to the node c_i and $e(t)$ incident to c_j . Denote $s_{ij} = t_i \dots t_{j-2} t_{j-1} t_{j-2} \dots t_i$ and $s_{ji} = t_{i-1} t_{i-2} \dots t_1 t_r \dots t_{j+1} t_j t_{j+1} \dots t_r t_1 \dots t_{i-1}$.

The following result shows that when Γ_X contains a circle with more than two nodes, the branching-circle relations in (vi) are a consequence of the branching relations in (v).

Lemma 4.12. *In the above setup, if the reflection set X satisfies all the o.b. relations together with the branching relations $st_{i-1}t_it_{i-1} = t_{i-1}t_it_{i-1}s$ and $tt_{j-1}t_jt_{j-1} = t_{j-1}t_jt_{j-1}t$ then*

$$(1) \quad ss_{ji}s \cdot ts_{ij}t = ts_{ij}t \cdot ss_{ji}s.$$

$$(2) \quad ss_{ji}s_{ij}ss_{ji}s_{ij} = s_{ji}s_{ij}ss_{ji}s_{ij}s.$$

$$(3) \quad ts_{ji}s_{ij}ts_{ji}s_{ij} = s_{ji}s_{ij}ts_{ji}s_{ij}t.$$

Proof. (1) holds $\iff t_j \dots t_r t_1 \dots t_{i-2} st_{i-1} st_{i-2} \dots t_1 t_r \dots t_j t_i \dots t_{j-2} tt_{j-1} tt_{j-2} \dots t_i$
 $= t_i \dots t_{j-2} tt_{j-1} tt_{j-2} \dots t_i t_j \dots t_r t_1 \dots t_{i-2} st_{i-1} st_{i-2} \dots t_1 t_r \dots t_j$
 $\iff t_j \dots t_r t_1 \dots t_{i-2} t_i \dots t_{j-2} st_{i-1} st_{j-1} tt_{i-2} \dots t_1 t_r \dots t_{j+1} t_j t_{j-2} \dots t_i$
 $= t_j \dots t_r t_1 \dots t_{i-2} t_i \dots t_{j-2} tt_{j-1} t st_{i-1} st_{i-2} \dots t_1 t_r \dots t_j t_{j-2} \dots t_i$
 $\iff st_{i-1} st_{j-1} t = tt_{j-1} t st_{i-1} s.$

If $j > i + 1$ and $(i, j) \neq (1, r)$, then s, t_{i-1} commute with t, t_{j-1} by the o.b. relations on X . So the last equation holds in this case. If $j = i + 1$ (resp., $(i, j) = (1, r)$), then the last equation becomes $st_{i-1} st_{i-1} t = tt_{i-1} t st_{i-1} s$ (resp., $st_r st_{r-1} t = tt_{r-1} t st_r s$). The last equation is again true by the o.b. and branching relation on X . So (1) follows.

Then (2) holds \iff

$$\begin{aligned} & t_{i-1} st_{i-1} t_{i-2} \dots t_1 t_r \dots t_j t_{j-1} t_{j-2} \dots t_{i+1} t_i t_{i+1} \dots t_{j-2} t_{j-1} st_j t_{j+1} \dots t_r t_1 \dots t_{i-2} t_{i-1} t_{i-2} \dots t_1 t_r \dots t_j t_{j-1} t_{j-2} \dots t_{i+1} \\ &= t_{i-2} \dots t_1 t_r \dots t_j t_{j-1} t_{j-2} \dots t_{i+1} t_i t_{i+1} \dots t_{j-2} t_{j-1} st_j t_{j+1} \dots t_r t_1 \dots t_{i-2} t_{i-1} t_{i-2} \dots t_1 t_r \dots t_j t_{j-1} t_{j-2} \dots t_{i+1} t_i st_i \\ &\iff t_j \dots t_r t_1 \dots t_{i-2} t_{i-1} st_{i-1} t_{i-2} \dots t_1 t_r \dots t_j t_i t_{i+1} \dots t_{j-2} t_{j-1} t_{j-2} \dots t_i st_j \dots t_r t_1 \dots t_{i-2} t_{i-1} t_{i-2} \dots t_1 t_r \dots t_j \\ &= t_{j-1} t_{j-2} \dots t_{i+1} t_i t_{i+1} \dots t_{j-2} t_{j-1} st_{i-1} t_{i-2} \dots t_1 t_r \dots t_{j+1} t_j t_{j+1} \dots t_r t_1 \dots t_{i-2} t_{i-1} t_{j-1} t_{j-2} \dots t_{i+1} t_i st_i t_{i+1} \dots t_{j-2} t_{j-1} \\ &\iff t_i \dots t_{j-2} t_j \dots t_r t_1 \dots t_{i-1} st_{i-1} \dots t_1 t_r \dots t_j t_{j-1} t_{j-2} \dots t_{i+1} t_i st_{i-1} t_{i-2} \dots t_1 t_r \dots t_{j+1} t_j t_{j+1} \dots t_r t_1 \dots t_{i-2} t_{i-1} \\ &= t_i \dots t_{j-2} t_{j-1} t_{j-2} \dots t_{i+1} t_i st_{i-1} t_{i-2} \dots t_1 t_r \dots t_{j+1} t_j t_{j-1} t_{j-2} \dots t_{i+1} t_i st_i t_{i+1} \dots t_{j-2} t_{j-1} t_{j+1} \dots t_r t_1 \dots t_{i-2} t_{i-1} \\ &\iff t_j \dots t_r t_1 \dots t_{i-2} t_{i-1} st_{i-1} t_{i-2} \dots t_1 t_r \dots t_{i+1} t_i st_{i-1} t_{i-2} \dots t_1 t_r \dots t_{j+1} t_j \\ &= t_{j-1} t_{j-2} \dots t_{i+1} t_i st_{i-1} t_{i-2} \dots t_1 t_r \dots t_{i+1} t_i st_i t_{i+1} \dots t_{j-2} t_{j-1} \\ &\iff t_i st_{i-1} t_{i-2} \dots t_1 t_r \dots t_{i+2} t_{i+1} t_{i+2} \dots t_r t_1 \dots t_{i-2} t_{i-1} st_i \\ &= st_{i-1} st_i t_{i-2} \dots t_1 t_r \dots t_{i+2} t_{i+1} t_{i+2} \dots t_r t_1 \dots t_{i-2} t_i st_{i-1} s \end{aligned}$$

$$\iff t_{i-2} \dots t_1 t_r \dots t_{i+2} t_{i+1} t_{i+2} \dots t_r t_1 \dots t_{i-2} = s t_{i-2} \dots t_1 t_r \dots t_{i+2} t_{i+1} t_{i+2} \dots t_r t_1 \dots t_{i-2} s.$$

The last equation is true since s commutes with all the other reflections occurring in the equation by the o.b. relations on X . So we have proved (2). Then (3) can be shown similarly. \square

Remark 4.13. Note that the conclusion of Lemma 4.12 still holds even when either $i + 1 = j$ or $j + 1 = i$ or both $i + 1 = j$, $j + 1 = i$ hold (the numbers are modulo r), where if $i + 1 = j$ (resp. $i = 1$ and $j = r$) then the element s_{ij} becomes t_i (resp. s_{ji} becomes t_j). When all the relations $i + 1 = j$, $i = 1$ and $j = r$ hold, we have $r = 2$. Thus the conclusion of Lemma 4.12 remains valid even when $r = 2$. However, in that case, it is not a consequence of the conditions $st_{i-1}t_it_{i-1} = t_{i-1}t_it_{i-1}s$ and $tt_{j-1}t_jt_{j-1} = t_{j-1}t_jt_{j-1}t$ since the latter no longer hold.

4.14. Let $X, X' \in \Sigma(m, 1, n)$ be such that X' is obtained from X by a terminal node operation with respect to two terminal nodes i, j along a path $a_1 = i, a_2, \dots, a_r = j$ for some $r > 2$. Hence $t_i = s(a_i, a_{i+1}; k_i) \in X$ for $1 \leq i < r$ and some integers k_i . Denote $t = t_1 t_2 \dots t_{r-2} t_{r-1} t_{r-2} \dots t_1$. We may assume $X' = (X \setminus \{t_1\}) \cup \{t\}$ without loss of generality.

Lemma 4.15. *In the above setup, if the reflection set X satisfies all the basic relations then so does the reflection set X' .*

Proof. We need only check all the basic relations on X' involving t . To do this, we need only check that

$$(1) \quad tt_{r-1}t = t_{r-1}tt_{r-1}.$$

(2) For any $s \in X' \setminus \{t, t_{r-1}\}$ with $e(s)$ incident to some a_h , $1 < h < r$, we have

$$(4.15.1) \quad st = ts.$$

(3) If $s = s(a_r; k)$ is in X then $stst = tsts$ and $s \cdot t_{r-1}tt_{r-1} = t_{r-1}tt_{r-1} \cdot s$.

(4) If $s = s(a_1; k)$ is in X then $stst = tsts$.

(1) is equivalent to the equation

$$t_{r-1} \dots t_2 t_1 t_2 \dots t_{r-2} t_{r-1} t_{r-2} \dots t_2 t_1 t_2 \dots t_{r-1} = t_{r-2} \dots t_2 t_1 t_2 \dots t_{r-2}.$$

The latter follows by the o.b. relations on t_1, \dots, t_{r-1} . For (2), s is in one of the following cases:

(i) $s = t_h$, $1 < h < r - 1$;

(ii) there is some node $a \neq a_1, \dots, a_r$ of Γ_X with $e(s) = \{a, a_h\}$ for some $1 < h < r$.

(iii) there is some h , $1 < h < h + 1 < r$ with $e(s) = \{a_h, a_{h+1}\}$ and $s \neq t_h$. Hence a_h and a_{h+1} are the nodes of a two-nodes circle in Γ_X .

(4.15.1) can be checked directly in case (i). Let $s = s(a, a_h; k)$ (resp. $s = s(a_h, a_{h+1}; k) \neq t_h$) be the reflection of X in case (ii) (resp. (iii)). Then in case (ii), a_h is a branching node of the graph Γ_X and so we have

$$(4.15.2) \quad s \cdot t_h t_{h-1} t_h = t_h t_{h-1} t_h \cdot s.$$

Now (4.15.1) holds $\iff t_{h-1} t_h \dots t_{r-2} t_{r-1} t_{r-2} \dots t_h t_{h-1} \cdot s = s \cdot t_{h-1} t_h \dots t_{r-2} t_{r-1} t_{r-2} \dots t_h t_{h-1}$
 $\iff t_{r-1} t_{r-2} \dots t_h t_{h-1} t_h \dots t_{r-1} \cdot s = s \cdot t_{r-1} t_{r-2} \dots t_h t_{h-1} t_h \dots t_{r-1} \iff (4.15.2)$ holds.

In case (iii), we have

$$(4.15.3) \quad t_{h-1} t_h t_{h-1} t_{h+1} s t_{h+1} = t_{h+1} s t_{h+1} t_{h-1} t_h t_{h-1}.$$

Now (4.15.1) holds

$$\iff t_{h-1} t_h t_{h+1} \dots t_{r-2} t_{r-1} t_{r-2} \dots t_{h+1} t_h t_{h-1} \cdot s = s \cdot t_{h-1} t_h t_{h+1} \dots t_{r-2} t_{r-1} t_{r-2} \dots t_{h+1} t_h t_{h-1}$$

$$\iff t_{r-1} \dots t_{h+1} t_h t_{h-1} t_h t_{h+1} \dots t_{r-1} \cdot s = s \cdot t_{r-1} \dots t_{h+1} t_h t_{h-1} t_h t_{h+1} \dots t_{r-1}$$

$$\iff (4.15.3) \text{ holds.}$$

Finally, (3) and (4) follows by the o.b. relations on s, t_1, \dots, t_{r-1} .

This shows our result. \square

Next result is the converse of Lemma 4.15.

Lemma 4.16. *Let $X, X' \in \Sigma(m, 1, n)$ be as in 4.14. If all the basic relations on X' hold then so do those on X .*

Proof. By Lemma 4.15, we need only consider the case where a_2 is not a terminal node of $\Gamma_{X'}$, i.e., there exists some $s \in X \setminus \{t_2\}$ with $e(s)$ incident to the node a_2 .

(i) First assume that the graph $\Gamma_{\{s, t_2\}}$ contains no two-nodes circle. Concerning the branching relations, it is sufficient, by Lemma 4.8, to show the relation $s \cdot t_2 t_1 t_2 = t_2 t_1 t_2 \cdot s$, that is to show:

$$(4.16.1) \quad t_2 s t_2 \cdot t_2 t_3 \dots t_{r-1} t t_{r-1} \dots t_2 = t_2 t_3 \dots t_{r-1} t t_{r-1} \dots t_2 \cdot t_2 s t_2.$$

This follows if the reflection s commutes with all of t_3, \dots, t_{r-1}, t . Otherwise, $e(s)$ is incident to the node a_i for some $3 < i \leq r-1$ (note that $e(s)$ is never incident to a_r since a_r is assumed a terminal node of $\Gamma_{X'}$). a_i is a branching node of $\Gamma_{X'}$ and we have the branching relation

$$(4.16.2) \quad s \cdot t_{i-1} t_i t_{i-1} = t_{i-1} t_i t_{i-1} \cdot s$$

on X' . relation (4.16.1) is equivalent to

$$(4.16.3) \quad s \cdot t_3 \dots t_{i-2} \cdot t t_{r-1} \dots t_i t_{i-1} t_i \dots t_{r-1} t \cdot t_{i-2} \dots t_3 \cdot s = t_3 \dots t_{i-2} \cdot t t_{r-1} \dots t_i t_{i-1} t_i \dots t_{r-1} t \cdot t_{i-2} \dots t_3.$$

Then (4.16.3) is an easy consequence of (4.16.2) together with some o.b. relations on X' .

If $s' = s(a_2; k)$ is in X , then we need also show the relations $s' t_1 s' t_1 = t_1 s' t_1 s'$ and $s' \cdot t_1 t_2 t_1 = t_1 t_2 t_1 \cdot s'$. Note $t_1 = t_2 t_3 \dots t_{r-1} t t_{r-1} \dots t_3 t_2$. So the relation $s' t_1 s' t_1 = t_1 s' t_1 s'$ is equivalent to

$$\begin{aligned} & s' \cdot t_2 t_3 \dots t_{r-1} t t_{r-1} \dots t_3 t_2 \cdot s' \cdot t_2 t_3 \dots t_{r-1} t t_{r-1} \dots t_3 t_2 \\ & = t_2 t_3 \dots t_{r-1} t t_{r-1} \dots t_3 t_2 \cdot s' \cdot t_2 t_3 \dots t_{r-1} t t_{r-1} \dots t_3 t_2 \cdot s'. \end{aligned}$$

The latter follows by the o.b. relations on $t, s', t_2, \dots, t_{r-1}$. The relation $s' \cdot t_1 t_2 t_1 = t_1 t_2 t_1 \cdot s'$ follows by the o.b. relations on $s', t, t_2, \dots, t_{r-1}$.

(ii) Next assume that the graph $\Gamma_{\{s, t_2\}}$ contains a two-nodes circle. Hence $r > 3$. We must show the branching-circle relation

$$(4.16.4) \quad t_1 t_2 t_1 \cdot t_3 s t_3 = t_3 s t_3 \cdot t_1 t_2 t_1.$$

This is equivalent to

$$(4.16.5) \quad t_3 s \cdot t_4 \dots t_{r-1} t t_{r-1} \dots t_4 \cdot s t_3 = t_3 \dots t_{r-1} t t_{r-1} \dots t_3$$

by applying some o.b. relations on X' . Now (4.16.5) follows since s commutes with t_4, \dots, t_{r-1}, t by the o.b. relations on X' . \square

Theorem 4.17. *For any $S \in \Sigma(m, 1, n)$, let P be the set of basic relations on S . Then (S, P) forms a presentation of the group $G(m, 1, n)$.*

Proof. Let $S \in \Sigma(m, 1, n)$ be given as in 4.1 with P the set of basic relations on S . Then (S, P) forms a presentation of $G(m, 1, n)$ by [1, Appendix 2]. Suppose that $X, X' \in \Sigma(m, 1, n)$ are as in 4.14 with P, P' the set of basic relations on X, X' respectively. Then by Lemmas 4.15 and 4.16, we see that (X, P) is a presentation of $G(m, 1, n)$ if and only if so is (X', P') . This implies our result by the fact that the set $\Sigma(m, 1, n)$ is transitive under terminal node operations (see 4.2). \square

4.18. Suppose that $X \in \Sigma(m, m, n)$ contains the reflections $t_h = s(c_h, c_{h+1}; k_h)$ (the subscripts are modulo r) for $1 \leq h \leq r$ and some integers k_h , where $r > 2$, and c_1, \dots, c_r are some nodes in Γ_X . Let $t = t_1 t_r t_1$ and let $X' = (X \setminus \{t_r\}) \cup \{t\}$. Then X' is obtained from X by a circle contraction.

Lemma 4.19. *In the above setup, the reflection set X satisfies all the basic relations if and only if so does the reflection set X' .*

Proof. First assume X satisfies all the basic relations. We want to show X' also satisfies all the basic relations. We need only check all the basic relations involving t . Note $e(t) = \{c_2, c_r\}$.

The order relation $t^2 = 1$ follows by the order relations $t_1^2 = 1 = t_r^2$ on X .

Let $s \in X' \setminus \{t\}$ be with $e(s)$ not incident to the edge $e(t)$. We must show $st = ts$. We see that $e(s)$ is incident to either both or none of $e(t_1), e(t_r)$. The result is obvious if $e(s)$ is incident to none of $e(t_1), e(t_r)$. In the case when $e(s)$ is incident to both of $e(t_1), e(t_r)$, we see that c_1 is a branching node of Γ_X to which the edges $e(t_1), e(t_r), e(s)$ incident. Then we have $ts = t_1 t_r t_1 s = s t_1 t_r t_1 = st$ by the branching relations on X .

Let $s \in X' \setminus \{t\}$ be with $e(s)$ incident to $e(t)$ at exactly one node in $\Gamma_{X'}$. We want to show $sts = tst$, i.e., $st_1 t_r t_1 s = t_1 t_r t_1 s t_1 t_r t_1$. This can be shown by the braid relations on X that either the relations $st_1 = t_1 s$, $st_r s = t_r s t_r$, or the relations $st_1 s = t_1 s t_1$, $st_r = t_r s$ hold. When $r = 3$, $e(t_2)$ and $e(t)$ form the two-nodes circle of $\Gamma_{X'}$. The circle relation $(t_2 t)^m = 1$ on X' is the same as the circle relation $(t_2 t_1 t_3 t_1)^m = 1$ on X .

So we have shown the o.b. relations on X' involving t .

Now we show the branching relations on X' involving t . If c_2 is a branching node in $\Gamma_{X'}$, then for any $s \in X' \setminus \{t_1, t\}$ with $e(s)$ incident to c_2 and not to c_r , we need show the relation $st_1 t t_1 = t_1 t t_1 s$. This follows by the braid relation $st_r = t_r s$ on X . If c_r is a branching node in $\Gamma_{X'}$, then for any $s \in X' \setminus \{t, t_{r-1}\}$ with $e(s)$ incident to c_r and not to c_2 , we need show the relation $st_{r-1} t t_{r-1} = t_{r-1} t t_{r-1} s$. This follows by the braid relation $t_1 s = s t_1$ and the branching relation $st_{r-1} t_r t_{r-1} = t_{r-1} t_r t_{r-1} s$ on X .

The circle relation $(ts_{2,r})^m = 1$ on X' is the same as the circle relation $(t_1 t_r t_1 s_{2,r})^m = 1$ on X .

It remains to show the branching-circle relations on X' involving t . By Lemma 4.12, we need only consider the case of $r = 3$. In this case, $e(t)$ and $e(t_2)$ form a two-nodes circle. If there exists some $u \in X' \setminus \{t, t_2\}$ with $e(u)$ incident to c_2 then the branching-circle relation $ut_2 t u t_2 t = t_2 t u t_2 t u$ on X' is the same as the branching-circle relation

$ut_2t_1t_3t_1ut_2t_1t_3t_1 = t_2t_1t_3t_1ut_2t_1t_3t_1u$ on X . Similarly for the case when there exists some $v \in X' \setminus \{t, t_2\}$ with $e(v)$ incident to c_r . If both of such u, v exist then the branching-circle relation $utuvvt_2v = vt_2vutu$ on X' is also the same as the branching-circle relation $ut_1t_3t_1uvt_2v = vt_2vut_1t_3t_1u$ on X .

Next assume that X' satisfies all the basic relations. We must show that so does the reflection set X . We need only show all the basic relations on X which involve the reflection $t_r = t_1tt_1$. Hence we have to show the following relations:

- (1) $t_r^2 = 1$;
- (2) $t_rs = st_r$ for $s \in X \setminus \{t_r\}$ with $e(s), e(t_r)$ having no common node;
- (3) $t_rst_r = st_rs$ for $s \in X \setminus \{t_r\}$ with $e(s), e(t_r)$ having exactly one common node;
- (4) $(t_r \cdot t_1t_2 \dots t_{r-2}t_{r-1}t_{r-2} \dots t_2t_1)^m = 1$;
- (5) $x \cdot t_1t_rt_1 = t_1t_rt_1 \cdot x$ if $x \in X$ is with $e(x)$ incident to the circle of Γ_X at the node c_1 ;
- (6) $y \cdot t_rt_{r-1}t_r = t_rt_{r-1}t_r \cdot y$ if $y \in X$ is with $e(y)$ incident to the circle of Γ_X at the node c_r .

The proof for the above relations is similar to what we did before and hence is left to the readers. Note that the branching-circle relations on X is a consequence of the branching relations on X by Lemma 4.12 and hence they need not be checked. \square

Theorem 4.20. *Let $S \in \Sigma(m, m, n)$ and let P be the set of all the basic relations on S . Then (S, P) forms a presentation of the group $G(m, m, n)$.*

Proof. Let $S \in \Sigma(m, m, n)$ be such that Γ_S is a string with a two-nodes circle at one end. [1, Appendix 2] tells us that (S, P) forms a presentation of $G(m, m, n)$.

By Lemma 2.14, any reflection set $X \in \Sigma(m, m, n)$ can be transformed to some $S' \in \Sigma(m, m, n)$ by a sequence of circle contractions followed by a sequence of terminal node operations, where $\Gamma_{S'}$ is a string with a two-nodes circle at one end. Hence S and S' are congruent. So our result follows by Lemmas 4.15, 4.16 and 4.19. \square

Remark 4.21. For any $S \in \Sigma(m, 1, n)$ (resp. $S \in \Sigma(m, m, n)$), we can get a presentation (S, P) of the group $G(m, 1, n)$ (resp. $G(m, m, n)$) by Theorem 4.17 (resp. Theorem 4.20). However, such a presentation is not essential in general (see 1.7 and Lemmas 4.8, 4.12). For example, removing any $n - 2$ (resp., $n - 1$) of the $n - 1$ (resp., n) order relations of type I in P for such a presentation (S, P) of the group $G(m, 1, n)$ (resp., $G(m, m, n)$), denote by P' the resulting relation set. Then (S, P') still form a presentation of the group. I conjecture that it will become essential after removing all the redundant order, branching, branching-circle and root-braid relations mentioned above and in Lemmas 4.8, 4.10, 4.12.

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