# On list 3-dynamic coloring of near-triangulations 

Ruijuan $\mathrm{Gu}^{\mathrm{a}}$, Seog-Jin Kim ${ }^{\mathrm{b}}$, Yulai $\mathrm{Ma}^{\mathrm{c}, *}$, Yongtang $\mathrm{Shi}^{\mathrm{c}}$<br>${ }^{a}$ Sino-European Institute of Aviation Engineering, Civil Aviation University of China, Tianjin 300300, China<br>${ }^{b}$ Department of Mathematics Educations, Konkuk University, Republic of Korea<br>${ }^{c}$ Center for Combinatorics and LPMC, Nankai University, Tianjin 300071, China


#### Abstract

An $r$-dynamic $k$-coloring of a graph $G$ is a proper $k$-coloring such that for any vertex $v$, there are at least $\min \left\{r, \operatorname{deg}_{G}(v)\right\}$ distinct colors in $N_{G}(v)$. The $r$-dynamic chromatic number $\chi_{r}^{d}(G)$ of a graph $G$ is the least $k$ such that there exists an $r$-dynamic $k$-coloring of $G$. The list $r$-dynamic chromatic number of a graph $G$ is denoted by $c h_{r}^{d}(G)$. Loeb et al. [10] showed that $\operatorname{ch}_{3}^{d}(G) \leq 10$ for every planar graph $G$, and there is a planar graph $G$ with $\chi_{3}^{d}(G)=7$.

In this paper, we study a special class of planar graphs which have better upper bounds of $\operatorname{ch}_{3}^{d}(G)$. We prove that $c h_{3}^{d}(G) \leq 6$ if $G$ is a planar graph which is a near-triangulation, where a near-triangulation is a planar graph whose bounded faces are all 3 -cycles. Keywords: list $r$-dynamic coloring, planar graphs, triangulation, near-triangulation


## 1. Introduction

Let $k$ be a positive integer. A proper $k$-coloring $\phi: V(G) \rightarrow\{1,2, \ldots, k\}$ of a graph $G$ is an assignment of colors to the vertices of $G$ so that any two adjacent vertices receive distinct colors. The chromatic number $\chi(G)$ of a graph $G$ is the least $k$ such that there exists a proper $k$-coloring of $G$. An $r$-dynamic $k$-coloring of a graph $G$ is a proper $k$ coloring $\phi$ such that for each vertex $v \in V(G)$, either the number of distinct colors in its neighborhood is at least $r$ or the colors in its neighborhood are all distinct, that is, $\left|\phi\left(N_{G}(v)\right)\right| \geq \min \left\{r, \operatorname{deg}_{G}(v)\right\}$. The $r$-dynamic chromatic number $\chi_{r}^{d}(G)$ of a graph $G$ is the least $k$ such that there exists an $r$-dynamic $k$-coloring of $G$.

A list assignment on a graph $G$ is a function $L$ that assigns each vertex $v$ a set $L(v)$ which is a list of available colors at $v$. For a list assignment $L$ of a graph $G$, we say $G$ is

[^0]$L$-colorable if there exists a proper coloring $\phi$ such that $\phi(v) \in L(v)$ for every $v \in V(G)$. A graph $G$ is said to be $k$-choosable if for any list assignment $L$ such that $|L(v)| \geq k$ for every vertex $v, G$ is $L$-colorable.

For a list assignment $L$ of $G$, we say that $G$ is $r$-dynamically $L$-colorable if there exists an $r$-dynamic coloring $\phi$ such that $\phi(v) \in L(v)$ for every $v \in V(G)$. A graph $G$ is $r$ dynamically $k$-choosable if for any list assignment $L$ with $|L(v)| \geq k$ for every vertex $v$, $G$ is $r$-dynamically $L$-colorable. The list $r$-dynamic chromatic number or the $r$-dynamic choice number $c_{r}^{d}(G)$ of a graph $G$ is the least $k$ such that $G$ is $r$-dynamically $k$-choosable.

An interesting property of dynamic coloring is as follows.

$$
\chi(G) \leq \chi_{2}^{d}(G) \leq \cdots \leq \chi_{\Delta}^{d}(G)=\chi\left(G^{2}\right)
$$

where $G^{2}$ is the square of the graph $G$.
The dynamic coloring was first introduced in $[8,11]$. On the other hand, Wegner [14] conjectured that if $G$ is a planar graph, then

$$
\chi_{\Delta}^{d}(G) \leq \begin{cases}7, & \text { if } \Delta(G)=3 ; \\ \Delta(G)+5, & \text { if } 4 \leq \Delta(G) \leq 7 ; \\ \left\lfloor\frac{3 \Delta(G)}{2}\right\rfloor+1, & \text { if } \Delta(G) \geq 8 .\end{cases}
$$

Lai et al. [12] posed a similar conjecture about dynamic coloring of planar graphs as follows.

Conjecture 1.1 Let $G$ be planar graph. Then

$$
\chi_{r}^{d}(G) \leq \begin{cases}r+3, & \text { if } 1 \leq r \leq 2 ; \\ r+5, & \text { if } 3 \leq r \leq 7 ; \\ \left\lfloor\frac{3 r}{2}\right\rfloor+1, & \text { if } r \geq 8 .\end{cases}
$$

Lai et al. [13] showed that conjecture 1.1 is true for planar graphs with girth at least 6. For the special case $r=2$, Kim et al. [6] proved that $\chi_{2}^{d}(G) \leq 4$ for every planar graph except $C_{5}$ and $c h_{2}^{d}(G) \leq 5$ for every planar graph. And it was shown in [10] that $c h_{3}^{d}(G) \leq 10$ if $G$ is a planar graph. Besides, some special classes of graphs are also investigated, such as sparse graphs [2], bipartite graphs [3], grids [4, 5], $K_{1,3}$-free graphs [9] and $K_{4}$-minor free graphs [12]. In terms of the maximum average degree, there is also a result published in [7].

Loeb et al. [10] showed $\operatorname{ch}_{3}^{d}(G) \leq 10$ if $G$ is a planar graph. On the other hand, there is a planar graph $F$ such that $\chi_{3}^{d}(F)=7$. So Loeb et al. [10] proposed the following problem.

Problem 1 ([10]) What is $\chi_{3}^{d}(G)$ if $G$ is a planar graph? And what is $c h_{3}^{d}(G)$ if $G$ is a planar graph?

Currently, we have the following bounds.

$$
\begin{equation*}
7 \leq \max \left\{\chi_{3}^{d}(G): G \text { is a planar graph }\right\} \leq 10 \tag{1}
\end{equation*}
$$

It is natural to consider a special class of planar graphs for Problem 1. Recently, Asayama et al. [1] showed that $\chi_{3}^{d}(G) \leq 5$ if $G$ is a triangulated planar graph, and the upper bound is sharp. But, we do not know yet whether $c h_{3}^{d}(G) \leq 5$ or not, if $G$ is a triangulated planar graph. So the following question is still open and it would be interested to answer.

Question 1 Is it true that $\operatorname{ch}_{3}^{d}(G) \leq 5$ if $G$ is a triangulated planar graph?
Since there is a gap (1) for the general case of planar graphs, it would be interesting to study list 3-dynamic chromatic number $c h_{3}^{d}(G)$ for a special class of planar graphs. In this paper, we consider a near-triangulation where a near-triangulation is a planar graph whose bounded faces are all 3 -cycles and outer face is bounded by a cycle. Note that a triangulated planar graph is a special case of a near-triangulation. First, we show the following theorem.

Theorem 1.2 If $G$ is a near-triangulation, then $\operatorname{ch}_{3}^{d}(G) \leq 6$.
And we obtain the following corollary.
Corollary 1.3 If $G$ is a triangulated planar graph, then $\operatorname{ch}_{3}^{d}(G) \leq 6$.
Let $W_{n}$ be the wheel with $n+1$ vertices such that $W_{n}$ is obtained from an $n$-cycle by adding a new vertex $u$ and joining $u$ and every vertex on the $n$-cycle. The following can be easily checked.

Proposition $1.4 \operatorname{ch}_{3}^{d}\left(W_{n}\right) \leq 6$ for every positive integer $n \geq 3$ and $c h_{3}^{d}\left(W_{5}\right)=6$.
Note that Proposition 1.4 and Theorem 1.2 imply that the upper bound of list 3dynamic chromatic number of near triangulations is tight. And Corollary 1.3 and [1] imply that

$$
5 \leq \max \left\{c h_{3}^{d}(G): G \text { is a triangulated planar graph }\right\} \leq 6
$$

## 2. Proof of Theorem 1.2

Suppose that Theorem 1.2 does not hold, and let $G$ be a minimal counterexample in terms of the number $\sigma(G)=|V(G)|+|E(G)|$ to Theorem 1.2. Let $C: v_{0} v_{1} \cdots v_{t-1} v_{0}$ in counter-clockwise order be the boundary of the outer face of a plane graph $G$. If $|V(G)| \leq 6$, then it is easy to obtain $c h_{3}^{d}(G) \leq 6$, a contradiction. Hence we have $|V(G)| \geq 7$.

First, we prove the following Claim.
Claim 1 For any $v \in V(C)$, we have that $d_{G}(v) \geq 4$.
Proof. Suppose that there is a vertex $v_{k} \in V(C)$ with $d_{G}\left(v_{k}\right) \leq 3$. Let $u_{0}, u_{1}, \ldots, u_{s-1}$ denote the neighbors of $v_{k}$ in counter-clockwise order such that $v_{k} u_{i} u_{i+1}$ is a 3 -face for each $i \in\{0,1, \ldots, s-2\}$. And let $u_{0}=v_{k+1}(k+1$ are computed by modulo $t)$.

If $d_{G}\left(v_{k}\right)=2$ or $d_{G}\left(v_{k}\right)=3$ with $u_{0} u_{2} \in E(G)$, then we remove $v_{k}$ from $G$ and call the resulting graph by $G^{\prime}$. If $d_{G}\left(v_{k}\right)=3$ and $u_{0} u_{2} \notin E(G)$, then we remove $v_{k}$ from $G$ and add the edge $u_{0} u_{2}$ in the outer face, and call the resulting graph by $G^{\prime}$. Clearly, for all cases above, $G^{\prime}$ is a near-triangulation.

Let $L^{\prime}(v)=L(v)$ for every $v \in V\left(G^{\prime}\right)$. Since $G$ is a minimal counterexample, $G^{\prime}$ has a 3 -dynamic $L^{\prime}$-coloring $\phi$.

If $d_{G}\left(v_{k}\right)=2$, then there exists a vertex $u_{0}^{\prime}$ such that $u_{0}^{\prime} \in\left(N_{G}\left(u_{0}\right) \cap N_{G}\left(u_{1}\right)\right) \backslash\left\{v_{k}\right\}$ since $|V(G)| \geq 7$. Then we color $v_{k}$ by a color $c \in L\left(v_{k}\right) \backslash\left\{\phi\left(u_{0}^{\prime}\right), \phi\left(u_{0}\right), \phi\left(u_{1}\right)\right\}$, and we obtain that $G$ has a 3 -dynamic coloring from the list assignment $L$, a contradiction.

If $d_{G}\left(v_{k}\right)=3$, then the vertices $u_{0}, u_{1}$ and $u_{2}$ receive distinct colors under the coloring $\phi$. Suppose $u_{0} u_{2} \in E(G)$. Then we color $v_{k}$ by a color $c \in L\left(v_{k}\right) \backslash\left\{\phi\left(u_{0}\right), \phi\left(u_{1}\right), \phi\left(u_{2}\right)\right\}$, and we obtain that $G$ has a 3-dynamic coloring from the list assignment $L$. This is a contradiction. Hence suppose that $u_{0} u_{2} \notin E(G)$. If there is a vertex $u_{i}$ for $i \in\{0,2\}$ such that $\phi\left(N_{G}\left(u_{i}\right)\right)$ has at most two different colors, then we must color $v_{k}$ by a color $c \in L\left(v_{k}\right) \backslash\left(\phi\left(N_{G}\left(v_{k}\right)\right) \cup \phi\left(N_{G}\left(u_{i}\right)\right)\right)$ so that vertex $u_{i}$ satisfies the conditions of 3-dynamic coloring. Then one can easily check that the number of forbidden colors at $v_{k}$ is at most 5 as follows.

Let $S$ be the set consisting of the forbidden colors at $v_{k}$. If $\left|\phi\left(N_{G}\left(u_{0}\right)\right)\right| \geq 3$ and $\left|\phi\left(N_{G}\left(u_{2}\right)\right)\right| \geq 3$, then $S=\left\{\phi\left(u_{0}\right), \phi\left(u_{1}\right), \phi\left(u_{2}\right)\right\}$. If $\left|\phi\left(N_{G}\left(u_{i}\right)\right)\right| \leq 2$ and $\left|\phi\left(N_{G}\left(u_{j}\right)\right)\right| \geq 3$ for $\{i, j\}=\{0,2\}$, then $S=\left\{\phi\left(u_{0}\right), \phi\left(u_{1}\right), \phi\left(u_{2}\right)\right\} \cup \phi\left(N_{G}\left(u_{i}\right)\right)$. If $\left|\phi\left(N_{G}\left(u_{0}\right)\right)\right| \leq 2$ and $\left|\phi\left(N_{G}\left(u_{2}\right)\right)\right| \leq 2$, then $S=\left\{\phi\left(u_{0}\right), \phi\left(u_{1}\right), \phi\left(u_{2}\right)\right\} \cup \phi\left(N_{G}\left(u_{0}\right)\right) \cup \phi\left(N_{G}\left(u_{2}\right)\right)$. Since $u_{1} \in N_{G}\left(u_{0}\right) \cap N_{G}\left(u_{2}\right)$, we can easily obtain $|S| \leq 5$ for all cases above.

Thus we can color $v_{k}$ by a color $c \in L\left(v_{k}\right) \backslash S$ so that $G$ has a 3 -dynamic coloring from the list assignment $L$, and it implies that $G$ is 3 -dynamically $L$-colorable. This is a contradiction, which completes the proof of Claim 1.

Next, we prove the following Claim.
Claim 2 For any $w \in V(G) \backslash V(C)$, we have that $d_{G}(w) \geq 6$.
Proof. Suppose that there is a vertex $w$ with $d_{G}(w) \leq 5$. Let $w_{0}, w_{1}, \ldots, w_{s-1}$ denote the neighbors of $w$ in counter-clockwise order.

Suppose $d_{G}(w)=3$. We remove $w$ from $G$ and call the resulting graph by $G^{\prime}$. Let $L^{\prime}(v)=L(v)$ for every $v \in V\left(G^{\prime}\right)$. Since $G$ is a minimal counterexample, $G^{\prime}$ has a 3dynamic $L^{\prime}$-coloring $\phi$. So, we can color $w$ by a color $c \in L(w) \backslash \phi\left(N_{G}(w)\right)$ so that $G$ has a 3-dynamic coloring from the list assignment $L$ since $|L(v)| \geq 6$ for each $v \in V(G)$, a contradiction.

Now we suppose $4 \leq d_{G}(w) \leq 5$. With Claim 1 and the preceding paragraph, we suppose $d_{G}(v) \geq 4$ for each $v \in V(G)$. Then we remove $w$ from $G$ and add edges in the face formed by $\left\{w_{0}, w_{1}, \ldots, w_{s-1}\right\}$ so that the resulting graph, denoted by $G^{\prime}$, is a near-triangulation. Let $L^{\prime}(v)=L(v)$ for every $v \in V\left(G^{\prime}\right)$. Since $G$ is a minimal counterexample, $G^{\prime}$ has a 3-dynamic $L^{\prime}$-coloring $\phi$. Clearly, we have that $\left|\phi\left(N_{G}(w)\right)\right| \geq 3$. If $N_{G}(w)=\left\{w_{0}, w_{1}, \ldots, w_{s-1}\right\}$ has all different colors in the coloring $\phi$, then we color $w$ by a color $c \in L(w) \backslash\left\{\phi\left(w_{i}\right): 0 \leq i \leq s-1\right\}$. Then this gives a 3 -dynamic coloring from its list assignment $L$, a contradiction.

Next, we consider the case when $N_{G}(w)=\left\{w_{0}, w_{1}, \ldots, w_{s-1}\right\}$ has less than $s$ colors. Let $S=\left\{w_{i} \in N_{G}(w) \mid \phi\left(w_{i-1}\right)=\phi\left(w_{i+1}\right)\right\}$. Since $4 \leq d_{G}(w) \leq 5$ and $\left|\phi\left(N_{G}(w)\right)\right|<s$, we have that $S \neq \emptyset$ in this case. Note that $G$ is a near-triangulation and $d_{G}(v) \geq 4$ for each $v \in V(G)$. So for each $w_{i} \in S$, we can select a vertex $w_{i}^{\prime}$ such that $w_{i}^{\prime} \in$ $\left(N_{G}\left(w_{i}\right) \cap\left(N_{G}\left(w_{i-1}\right) \cup N_{G}\left(w_{i+1}\right)\right)\right) \backslash\left\{w, w_{i-1}, w_{i+1}\right\}(i-1$ and $i+1$ are computed by modulo $s$ ). Clearly, $\phi\left(w_{i}^{\prime}\right) \neq \phi\left(w_{i-1}\right)$ or $\phi\left(w_{i}^{\prime}\right) \neq \phi\left(w_{i+1}\right)$ since $\phi$ is a proper coloring. Now let $S^{\prime}=\left\{w_{i}^{\prime} \mid w_{i} \in S\right\}$.

Since $S \neq \emptyset$, we have that $\left|S^{\prime}\right| \geq 1$. And it is easy to check that $|S| \leq 2$ and $\left|S^{\prime}\right| \leq 2$ since $4 \leq d_{G}(w) \leq 5$ and $\left|\phi\left(N_{G}(w)\right)\right| \geq 3$. Moreover, if $|S|=1$, then $\left|S^{\prime}\right|=1$ and $\left|\phi\left(N_{G}(w)\right)\right| \leq 4$. If $|S|=2$, then $\left|S^{\prime}\right| \leq 2$ and $\left|\phi\left(N_{G}(w)\right)\right| \leq 3$. So for both cases above, we obtain that $L(w) \backslash\left(\left\{\phi\left(w_{i}^{\prime}\right): w_{i}^{\prime} \in S^{\prime}\right\} \cup \phi\left(N_{G}(w)\right) \neq \emptyset\right.$ since $|L(w)| \geq 6$. Then we color $w$ by a color $c \in L(w) \backslash\left(\left\{\phi\left(w_{i}^{\prime}\right): w_{i}^{\prime} \in S^{\prime}\right\} \cup \phi\left(N_{G}(w)\right)\right.$. Clearly, there are at least 3
distinct colors in $N_{G}(w)$ and at least 3 distinct colors in $N_{G}\left(w_{i}\right)$ for each $w_{i} \in N_{G}(w) \backslash S$. For each vertex $w_{i} \in S$, we have that $\phi\left(w_{i-1}\right)=\phi\left(w_{i+1}\right)$ and then $\phi\left(w_{i}^{\prime}\right) \neq \phi\left(w_{i-1}\right) \neq c$. So each $w_{i} \in S$ also satisfies the conditions of 3 -dynamic coloring. Thus we obtain that $G$ has a 3-dynamic coloring from the list assignment $L$, which is a contradiction since $G$ is a counterexample. This completes the proof of Claim 2.

Let $k$ be the number of vertices in $V(G) \backslash V(C)$. Then $n(G)=t+k$ since $|V(C)|=t$. Now from Claim 1 and Claim 2, we have

$$
\begin{equation*}
2 e(G)=\sum_{v \in V(G)} d_{G}(v)=\sum_{v \in V(C)} d_{G}(v)+\sum_{v \in V(G) \backslash V(C)} d_{G}(v) \geq 4 t+6 k . \tag{2}
\end{equation*}
$$

And since $G$ is a near-triangulation, we have

$$
\begin{equation*}
e(G)=3 n(G)-6-(|V(C)|-3)=3 n(G)-t-3=2 t+3 k-3 \tag{3}
\end{equation*}
$$

So, by (2) and (3)

$$
4 t+6 k-6=2 e(G) \geq 4 t+6 k \Longrightarrow-6 \geq 0
$$

which is a contradiction. This completes the proof of Theorem 1.2.
Acknowledgments. The authors are grateful to the editor and two referees for their valuable comments and constructive suggestions, and thankful to one of two referees who pointed out that Problem 1 was proposed by Loeb et al. [10]. Y. Ma and Y. Shi were partially supported by the National Natural Science Foundation of China (No. 11811540390). S.-J. Kim's work was supported by the National Research Foundation of Korea (NRF) grant funded by the Korea government (MSIT) (NRF-2018R1A2B6003412).

## References

[1] Y. Asayama, Y. Kawasaki, S.-J. Kim, A. Nakamoto, K. Ozeki, 3-dynamic coloring of planar triangulations, Discrete Math. 341 (2018) 2988-2994.
[2] J. Cheng, H.-J. Lai, K.J. Lorenzen, R. Luo, J.C. Thompson, C.-Q. Zhang, r-hued coloring of sparse graphs, Discrete Appl. Math. 237 (2018) 75-81.
[3] L. Esperet, Dynamic list coloring of bipartite graphs, Discrete Appl. Math. 158 (2010) 1963-1965.
[4] S. Jahanbekam, J. Kim, S. O, D.B. West, On r-dynamic coloring of graphs, Discrete Appl. Math. 206 (2016) 65-72.
[5] R. Kang, T. Müller, D.B. West, On $r$-dynamic coloring of grids, Discrete Appl. Math. 186 (2015) 286-290.
[6] S.-J. Kim, S. Lee, W. Park, Dynamic coloring and list dynamic coloring of planar graphs, Discrete Appl. Math. 161 (2013) 2207-2212.
[7] S.-J. Kim, B. Park, List 3-dynamic coloring of graphs with small maximum average degree, Discrete Math. 341 (5) (2018) 1406-1418.
[8] H.-J. Lai, B. Montgomery, H. Poon, Upper bounds of dynamic chromatic number, Ars Combin. 68 (2003) 193-201.
[9] H. Li, H.-J. Lai, 3-dynamic coloring and list 3-dynamic coloring of $K_{1,3}$-free graphs, Discrete Appl. Math. 222 (2017) 166-171.
[10] S. Loeb, T. Mahoney, B. Reiniger, J. Wise, Dynamic coloring parameters for graphs with given genus, Discrete Appl. Math. 235 (2018) 129-141.
[11] B. Montgomery, Ph.D. Dissertation, West Virginia University, 2001.
[12] H. Song, S. Fan, Y. Chen, L. Sun, H.-J. Lai, On $r$-hued coloring of $K_{4}$-minor free graphs, Discrete Math. 315 (2014) 47-52.
[13] H. Song, H.-J. Lai, J.-L. Wu, On $r$-hued coloring of planar graphs with girth at least 6, Discrete Appl. Math. 198 (2016) 251-263.
[14] G. Wegner, Graphs with given diameter and a coloring problem, Technical Report, University of Dortmund, 1977.


[^0]:    *Corresponding author
    Email addresses: millet90@163.com (Ruijuan Gu), skim12@konkuk.ac.kr (Seog-Jin Kim), ylma92@163.com (Yulai Ma), shi@nankai.edu. cn (Yongtang Shi)

