

Extremal Theta-free planar graphs

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Abstract

Given a family \mathcal{F} , a graph is \mathcal{F} -free if it does not contain any graph in \mathcal{F} as a subgraph. We continue to study the topic of “extremal” planar graphs initiated by Dowden [J. Graph Theory 83 (2016) 213–230], that is, how many edges can an \mathcal{F} -free planar graph on n vertices have? We define $ex_{\mathcal{P}}(n, \mathcal{F})$ to be the maximum number of edges in an \mathcal{F} -free planar graph on n vertices. Dowden obtained the tight bounds $ex_{\mathcal{P}}(n, C_4) \leq 15(n-2)/7$ for all $n \geq 4$ and $ex_{\mathcal{P}}(n, C_5) \leq (12n-33)/5$ for all $n \geq 11$. In this paper, we continue to promote the idea of determining $ex_{\mathcal{P}}(n, \mathcal{F})$ for certain classes \mathcal{F} . Let Θ_k denote the family of Theta graphs on $k \geq 4$ vertices, that is, graphs obtained from a cycle C_k by adding an additional edge joining two non-consecutive vertices. The study of $ex_{\mathcal{P}}(n, \Theta_4)$ was suggested by Dowden. We show that $ex_{\mathcal{P}}(n, \Theta_4) \leq 12(n-2)/5$ for all $n \geq 4$, $ex_{\mathcal{P}}(n, \Theta_5) \leq 5(n-2)/2$ for all $n \geq 5$, and then demonstrate that these bounds are tight, in the sense that there are infinitely many values of n for which they are attained exactly. We also prove that $ex_{\mathcal{P}}(n, C_6) \leq ex_{\mathcal{P}}(n, \Theta_6) \leq 18(n-2)/7$ for all $n \geq 6$.

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1 Introduction

All graphs considered in this paper are finite and simple. We use P_k and C_k to denote the path and cycle on k vertices, respectively. Let \mathcal{F} be a family of graphs. A graph is \mathcal{F} -free if it does not contain any graph in \mathcal{F} as a subgraph. When $\mathcal{F} = \{F\}$ we write F -free. One of the best known results in extremal graph theory is Turán’s Theorem [12], which gives the maximum number of edges that a K_k -free graph on n vertices can have. The celebrated Erdős-Stone Theorem [4] then extends this to the case when K_k is replaced by an arbitrary graph H , showing that the maximum number of edges possible is $(1 + o(1)) \left(\frac{\chi(H)-2}{\chi(H)-1} \right) n$, where $\chi(H)$ denotes the chromatic number of H . This latter result has been called the “fundamental theorem of extremal graph theory” [1]. Turán-type problems when host graphs are hypergraphs are notoriously difficult. A large quantity

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of work in this area has been carried out in determining the maximum number of edges in a k -uniform hypergraph on n vertices without containing k -uniform linear paths and cycles (see, for example, [6, 7, 10]). Surveys on Turán-type problems of graphs and hypergraphs can be found in [5] and [9].

Recently, Dowden [3] initiated the study of Turán-type problems when host graphs are planar graphs, i.e., how many edges can an \mathcal{F} -free planar graph on n vertices have? The *planar Turán number of \mathcal{F}* , denoted $ex_{\mathcal{P}}(n, \mathcal{F})$, is the maximum number of edges in an \mathcal{F} -free planar graph on n vertices. When $\mathcal{F} = \{F\}$ we write $ex_{\mathcal{P}}(n, F)$. Dowden [3] observed that it is straightforward to determine the exact values of $ex_{\mathcal{P}}(n, H)$ when H is a complete graph or non-planar graph; he also obtained the tight bounds $ex_{\mathcal{P}}(n, C_4) \leq 15(n-2)/7$ for all $n \geq 4$ and $ex_{\mathcal{P}}(n, C_5) \leq (12n-33)/5$ for all $n \geq 11$. Recently, Lan, Shi and Song observed in [11] that planar Turán numbers are closely related to planar anti-Ramsey numbers. The *planar anti-Ramsey number of \mathcal{F}* , denoted $ar_{\mathcal{P}}(n, \mathcal{F})$, is the maximum number of colors in an edge-coloring of a plane triangulation T (which is not \mathcal{F} -free) on n vertices such that T contains no rainbow copy of any $F \in \mathcal{F}$. When $\mathcal{F} = \{F\}$ we write $ar_{\mathcal{P}}(n, F)$. The study of planar anti-Ramsey numbers was initiated by Horňák, Jendrol', Schiermeyer and Soták [8] (under the name of rainbow numbers). The following result is observed in [11].

Proposition 1.1 ([11]) *Given a planar graph H and a positive integer $n \geq |H|$,*

$$1 + ex_{\mathcal{P}}(n, \mathcal{H}) \leq ar_{\mathcal{P}}(n, H) \leq ex_{\mathcal{P}}(n, H),$$

where $\mathcal{H} = \{H - e : e \in E(H)\}$.

In this paper, we continue to promote the idea of determining $ex_{\mathcal{P}}(n, \mathcal{F})$ for certain classes \mathcal{F} . This paper focuses on the family of Theta graphs, where a graph on at least 4 vertices is a *Theta graph* if it can be obtained from a cycle by adding an additional edge joining two non-consecutive vertices. For integer $k \geq 4$, let Θ_k be the family of non-isomorphic Theta graphs on k vertices. Note that the only graph in Θ_4 is isomorphic to K_4 minus one edge, and Θ_5 has only one graph. By abusing notation, we also use Θ_4 and Θ_5 to denote the only graph in Θ_4 and Θ_5 , respectively. Note that the study of $ex_{\mathcal{P}}(n, \Theta_4)$ was suggested by Dowden [3]. We need to introduce more notation. For a graph G , we will use $V(G)$ to denote the vertex set, $E(G)$ the edge set, $|G|$ the number of vertices, $e(G)$ the number of edges, $\delta(G)$ the minimum degree and \overline{G} the complement of G . For a vertex $x \in V(G)$, we will use $N_G(x)$ to denote the set of vertices in G which are adjacent to x . We define $d_G(x) = |N_G(x)|$. Given vertex sets $A, B \subseteq V(G)$, the subgraph of G induced on A , denoted $G[A]$, is the graph with vertex set A and edge set $\{xy \in E(G) : x, y \in A\}$. We denote by $B \setminus A$ the set $B - A$ and $G \setminus A$ the subgraph of G induced on $V(G) \setminus A$, respectively. We say that A is *complete to* (resp. *anti-complete to*) B if for every $a \in A$ and every $b \in B$, $ab \in E(G)$ (resp. $ab \notin E(G)$).

If $A = \{a\}$, we simply say a is complete to (resp. anti-complete to) B , and write $B \setminus a$ and $G \setminus a$, respectively. The *join* $G + H$ (resp. *union* $G \cup H$) of two vertex disjoint graphs G and H is the graph having vertex set $V(G) \cup V(H)$ and edge set $E(G) \cup E(H) \cup \{xy \mid x \in V(G), y \in V(H)\}$. (resp. $E(G) \cup E(H)$). For a positive integer t , we use tH to denote disjoint union of t copies of a graph H . Given two isomorphic graphs G and H , we may (with a slight but common abuse of notation) write $G = H$. A graph H is a *spanning subgraph* of a graph G if H is a subgraph of G with $V(H) = V(G)$. For any positive integer k , let $[k] := \{1, 2, \dots, k\}$.

We state and prove our main results in Section 2.

2 Planar Turán number of Theta graphs

In this section, using the method developed in [3], we study planar Turán numbers of Θ_k when $k \in \{4, 5, 6\}$. The study of $ex_{\mathcal{P}}(n, \Theta_4)$ was suggested by Dowden [3]. Our technique relies heavily on Euler's formula. We need to introduce more notation that shall be used in this section only.

An \mathcal{F} -free planar graph G on n vertices with the largest possible number of edges is called *extremal for n and \mathcal{F}* . If $\mathcal{F} = \{F\}$, then we simply say G is extremal for n and F . Given a plane graph G and integers $i, j \geq 3$, an *i -face* in G is a face of size i ; and let: $E_{i,j}$ denote the set of edges in G that each belong to one i -face and one j -face (and belong to two i -faces when $i = j$); E_i denote the set of edges in G that each belong to at least one i -face; and f_i denote the number of i -faces in G . Let $e_{i,j} := |E_{i,j}|$, $e_i := |E_i|$, and $f := \sum_i f_i$. Given three positive integers a, b and c , we use $a \equiv b \pmod{c}$ to denote a and b have the same remainder when divided by c . We will make use of the following observation.

Observation 2.1 *Let G be a plane graph on $n \geq 3$ vertices with $e(G) \geq 2$. For all $i \geq 3$,*

- (a) $e_{i,i} \leq e_i \leq e(G)$,
- (b) $if_i = e_i + e_{i,i}$,
- (c) $\sum_{i \geq 3} e_i - \sum_{3 \leq i < j} e_{i,j} = e(G)$, and
- (d) *every face in G is bounded by a cycle if G is 2-connected.*

We begin with $\mathcal{F} = \Theta_4$ and prove that $ex_{\mathcal{P}}(n, \Theta_4) \leq 12(n-2)/5$ for all $n \geq 4$ and then demonstrate that this bound is tight, in the sense that there are infinitely many values of n for which it is attained exactly.

Theorem 2.2 *$ex_{\mathcal{P}}(n, \Theta_4) \leq 12(n-2)/5$ for all $n \geq 4$, with equality when $n \equiv 12 \pmod{20}$.*

Proof. Let G be a Θ_4 -free plane graph on $n \geq 4$ vertices. We shall proceed the proof by induction on n . The statement is trivially true when $n = 4$ because any Θ_4 -free plane graph on

four vertices has at most four edges. So we may assume that $n \geq 5$. Next assume that there exists a vertex $u \in V(G)$ with $d_G(u) \leq 2$. By the induction hypothesis, $e(G \setminus u) \leq 12(n - 3)/5$ and so $e(G) = e(G \setminus u) + d_G(u) \leq 12(n - 3)/5 + 2 < 12(n - 2)/5$, as desired. So we may assume that $\delta(G) \geq 3$. Then each component of G has at least five vertices because G is Θ_4 -free. By the induction hypothesis, we may further assume that G is connected. Then G has no face of size at most two because G is simple. Hence

$$2e(G) = \sum_{i \geq 3} i f_i \geq 3f_3 + 4 \sum_{i \geq 4} f_i = 3f_3 + 4(f - f_3) = 4f - f_3, \tag{1}$$

which implies that $f \leq (2e(G) + f_3)/4$. Note that $E_{3,3} = \emptyset$ else G would contain Θ_4 as a subgraph, a contradiction. Thus $e_3 = 3f_3$ by Observation 2.1(b). This, together with $e_3 \leq e(G)$ and $f \leq (2e(G) + f_3)/4$, implies that $f \leq 7e(G)/12$. By Euler's formula, $n - 2 = e(G) - f \geq 5e(G)/12$. Hence $e(G) \leq 12(n - 2)/5$, as desired.

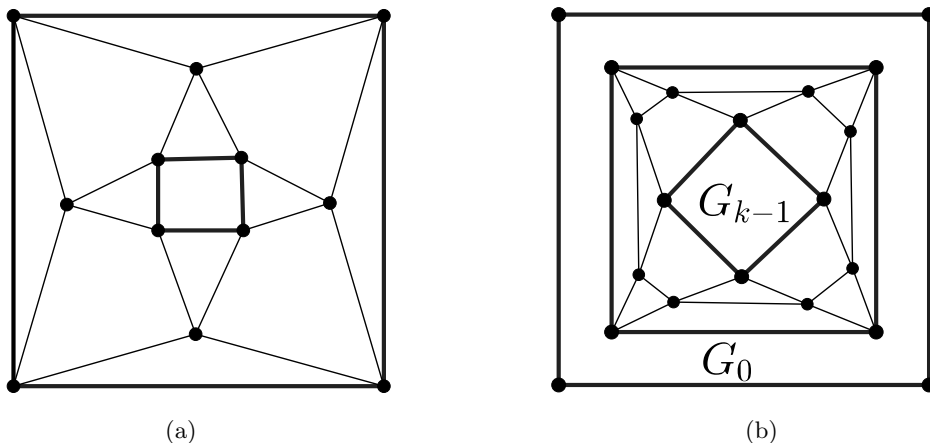


Figure 1: Construction of G_k .

From the proof above, we see that equality in $e(G) \leq 12(n - 2)/5$ is achieved for n if and only if equalities hold both in (1) and in $e_3 \leq e(G)$. This implies that $e(G) = 12(n - 2)/5$ for n if and only if G is a connected Θ_4 -free plane graph on n vertices such that each edge in G belongs to one 3-face and one 4-face. We next construct such an extremal graph for n and Θ_4 . Let $n = 20k + 12$ for some integer $k \geq 0$. Let G_0 be the graph depicted in Figure 1(a), we then construct G_k on n vertices recursively for all $k \geq 1$ via the illustration given in Figure 1(b): the entire graph G_{k-1} is placed into the center quadrangle of Figure 1(b), and the entire G_0 is then placed between the two given bold quadrangles of Figure 1(b) (in such a way that these are identified with the bold quadrangles of Figure 1(a)). One can check that G_k is Θ_4 -free with $n = 20k + 12$ vertices and $12(n - 2)/5$ edges for all $k \geq 0$. ■

We next prove that $ex_{\mathcal{P}}(n, \Theta_5) \leq 5(n-2)/2$ and then demonstrate that this bound is tight, in the sense that there are infinitely many values of n for which it is attained exactly.

Theorem 2.3 $ex_{\mathcal{P}}(n, \Theta_5) \leq 5(n-2)/2$ for all $n \geq 5$, with equality when $n \equiv 50 \pmod{120}$.

Proof. Let G be a Θ_5 -free plane graph on $n \geq 5$ vertices. We show by induction on n that $e(G) \leq 5(n-2)/2$. The statement is trivially true when $n = 5$ because any Θ_5 -free plane graph on five vertices has at most seven edges. So we may assume that $n \geq 6$. Next assume that there exists a vertex $u \in V(G)$ with $d_G(u) \leq 2$. By the induction hypothesis, $e(G \setminus u) \leq 5(n-3)/2$ and so $e(G) = e(G \setminus u) + d_G(u) \leq 5(n-3)/2 + 2 < 5(n-2)/2$, as desired. So we may assume that $\delta(G) \geq 3$. Assume next that G is disconnected. Let $G_1, \dots, G_s, G_{s+1}, \dots, G_{s+t}$ be all components of G such that $|G_1| = \dots = |G_s| = 4$ and $5 \leq |G_{s+1}| \leq \dots \leq |G_{s+t}|$, where $s \geq 0$ and $t \geq 0$ are integers with $s+t \geq 2$ and $4s + |G_{s+1}| + \dots + |G_{s+t}| = n$. Then $e(G_i) = 6$ for all $i \in [s]$ because $\delta(G) \geq 3$, and $e(G_j) \leq 5(|G_j| - 2)/2$ for all $j \in \{s+1, \dots, s+t\}$ by the induction hypothesis. Therefore,

$$\begin{aligned} e(G) &\leq 6s + \frac{5(|G_{s+1}| + \dots + |G_{s+t}| - 2t)}{2} \\ &= \frac{5(n-2)}{2} - \frac{(8(s+t) + 2t - 10)}{2} < \frac{5(n-2)}{2}, \end{aligned}$$

as desired. So we may further assume that G is connected.

Since G is a connected plane graph on $n \geq 6$ vertices, we see that G has no face of size at most two. Hence

$$2e(G) = 3f_3 + 4f_4 + \sum_{i \geq 5} if_i \geq 3f_3 + 4f_4 + 5(f - f_3 - f_4) = 5f - 2f_3 - f_4, \quad (2)$$

which implies that $f \leq (2e(G) + 2f_3 + f_4)/5$. Note that no 3-face in G has its three edges in $E_{3,3}$ because G is Θ_5 -free and $n \geq 6$. It follows that $e_{3,3} \leq f_3$. By Observation 2.1(b),

$$3f_3 = e_3 + e_{3,3} \leq e_3 + f_3 \text{ and so } f_3 \leq e_3/2. \quad (3)$$

It is worth noting that a 4-face and a 3-face in G cannot have exactly one edge in common, else G would contain Θ_5 as a subgraph. Since $\delta(G) \geq 3$, we see that a 4-face and a 3-face in G cannot have exactly two edges in common. Hence, every 4-face and every 3-face in G have no edge in common and so $E_{3,4} = \emptyset$. Thus, $e_3 + e_4 \leq e(G)$. By Observation 2.1(a,b), $e_{4,4} \leq e_4$ and $4f_4 = e_4 + e_{4,4}$. It follows that

$$4f_4 \leq 2e_4 \leq 2(e(G) - e_3) \text{ and so } f_4 \leq (e(G) - e_3)/2. \quad (4)$$

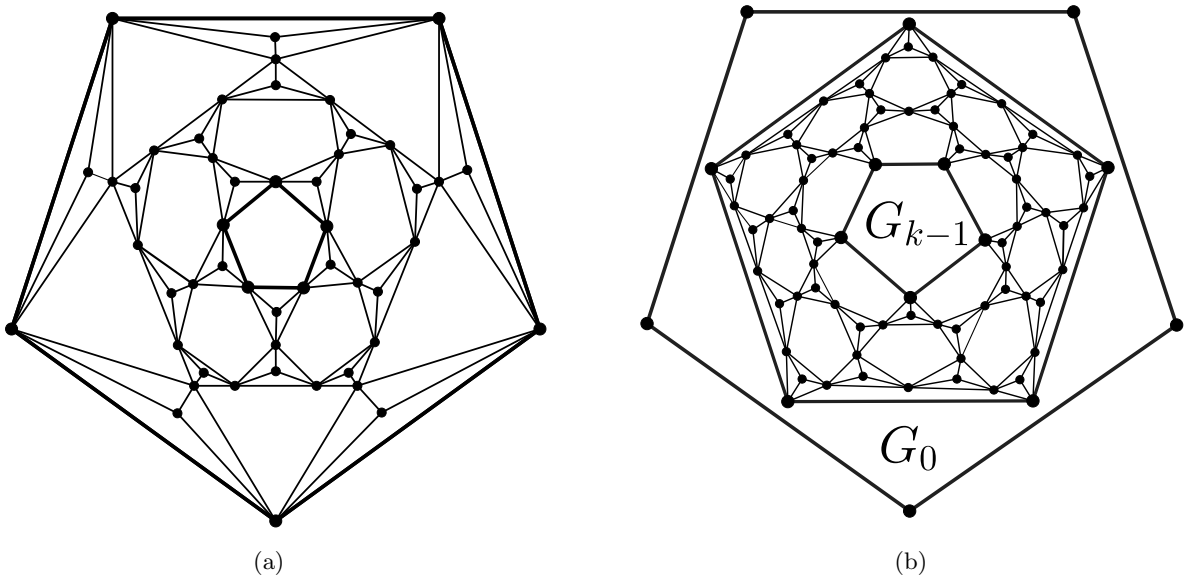


Figure 2: Construction of G_k .

Now with the last inequalities in (3) and (4), and the fact that $f \leq (2e(G) + 2f_3 + f_4)/5$ and $e_3 \leq e(G)$, we obtain $f \leq 3e(G)/5$. By Euler's formula, $n - 2 = e(G) - f \geq 2e(G)/5$. Hence $e(G) \leq 5(n - 2)/2$, as desired.

From the proof above, we see that equality in $e(G) \leq 5(n - 2)/2$ is achieved for n if and only if equalities hold in (2), (3) and (4) and in $e_3 \leq e(G)$. This implies that $e(G) = 5(n - 2)/2$ for n if and only if G is a connected Θ_5 -free plane graph on n vertices satisfying: each 3-face in G has exactly two edges in $E_{3,3}$; each edge in G belongs to either one 3-face and one 5-face or two 3-faces. We next construct such an extremal plane graph for n and Θ_5 . Let $n = 120k + 50$ for some integer $k \geq 0$. Let G_0 be the graph depicted in Figure 2(a), we then construct G_k of order n recursively for all $k \geq 1$ via the illustration given in Figure 2(b): the entire graph G_{k-1} is placed into the center pentagon of Figure 2(b), and the entire G_0 is then placed between the two given bold pentagons of Figure 2(b) (in such a way that these are identified with the bold pentagons of Figure 2(a)). One can check that G_k is Θ_5 -free with $n = 120k + 50$ vertices and $5(n - 2)/2$ edges for all $k \geq 0$. ■

Finally, we prove an upper bound for $ex_p(n, \Theta_6)$ in Theorem 2.4. Figure 3 illustrates all possible graphs for which equality in Theorem 2.4 is attained when $n = 9$. However, we shall see in Corollary 2.5 that equality is not possible for all $n \geq 10$.

Theorem 2.4 $ex_p(n, \Theta_6) \leq 18(n - 2)/7$ for all $n \geq 6$, with equality when $n = 9$.

Proof. Let G be an extremal plane graph for Θ_6 and $n \geq 6$. We shall prove that $e(G) \leq 18(n - 2)/7$ by induction on n . When $n = 6$, we show that $e(G) \leq 10$. Suppose that $e(G) \geq 11$.

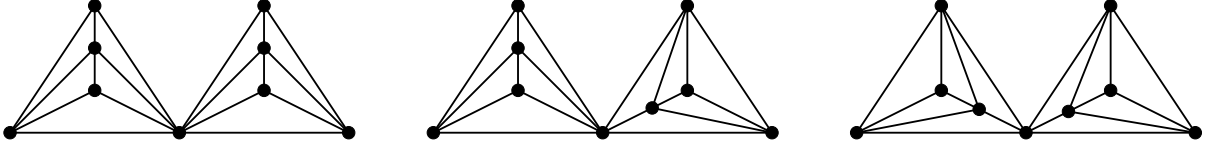


Figure 3: All possible graphs achieving equality in Theorem 2.4 and Corollary 2.5 when $n = 9$.

Then G is isomorphic to either a plane triangulation on six vertices or a plane triangulation on six vertices with one edge removed. Note that all plane triangulations on 6 vertices are depicted in Figure 4. It is easy to check that G has a Hamiltonian cycle and so G contains a graph in Θ_6 as subgraph, a contradiction. Hence, $e(G) \leq 10 < 18(n - 2)/7$ when $n = 6$. So we may assume that $n \geq 7$. Next assume that there exists a vertex $u \in V(G)$ with $d_G(u) \leq 2$. By the induction hypothesis, $e(G \setminus u) \leq 18(n - 3)/7$ and so $e(G) = e(G \setminus u) + d_G(u) \leq 18(n - 3)/7 + 2 < 18(n - 2)/7$, as desired. So we may assume that $\delta(G) \geq 3$. Assume next that G is disconnected. Then each component of G has exactly four, five or at least six vertices because $\delta(G) \geq 3$. Let $G_1, \dots, G_r, G_{r+1}, \dots, G_{r+s}, G_{r+s+1}, \dots, G_{r+s+t}$ be all components of G such that

$$|G_1| = \dots = |G_r| = 4, |G_{r+1}| = \dots = |G_{r+s}| = 5, \text{ and } 6 \leq |G_{r+s+1}| \leq \dots \leq |G_{r+s+t}|,$$

where $r, s, t \geq 0$ are integers with $r + s + t \geq 2$ and $4r + 5s + |G_{r+s+1}| + \dots + |G_{r+s+t}| = n$. Since G is an extremal plane graph for Θ_6 , we see that $e(G_i) = 6$ for all $i \in [r]$ and $e(G_j) = 9$ for all $j \in \{r + 1, \dots, r + s\}$. By the induction hypothesis, $e(G_k) \leq 18(|G_k| - 2)/7$ for all $k \in \{r + s + 1, \dots, r + s + t\}$. Therefore,

$$\begin{aligned} e(G) &\leq 6r + 9s + \frac{18(|G_{r+s+1}| + \dots + |G_{r+s+t}| - 2t)}{7} \\ &= \frac{18(n - 2)}{7} - \frac{(27(r + s + t) + 3r + 9t - 36)}{7} \\ &< \frac{18(n - 2)}{7}, \end{aligned}$$

as desired. So we may assume that G is connected.

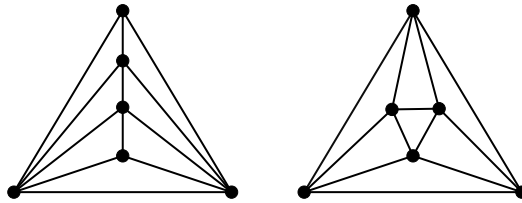


Figure 4: All plane triangulations on 6 vertices.

Next assume that G contains a cut-vertex, say u . Let H be a smallest component of $G \setminus u$, and let $G_1 := G[V(H) \cup \{u\}]$ and $G_2 := G \setminus V(H)$. Then $|G_1| \leq |G_2|$ and $|G_1| + |G_2| = n + 1$. Since $\delta(G) \geq 3$, we see that $4 \leq |G_1| \leq |G_2|$. Assume first that $|G_2| \leq 5$. Then $e(G_i) \leq 3|G_i| - 6$ for all $i \in \{1, 2\}$. Hence, $e(G) = e(G_1) + e(G_2) \leq 3(|G_1| + |G_2|) - 12 = 3n - 9 \leq 18(n - 2)/7$ because $n \leq 9$, with equality when both G_1 and G_2 are isomorphic to K_5 minus one edge, and so G is isomorphic to the graphs depicted in Figure 3. Assume next that $|G_2| \geq 6$. Then $e(G_2) \leq 18(|G_2| - 2)/7$ by the induction hypothesis. Note that $e(G_1) \leq 3|G_1| - 6$ when $|G_1| \leq 5$ and $e(G_1) \leq 18(|G_1| - 2)/7$ when $|G_1| \geq 6$ by the induction hypothesis. Therefore, when $|G_1| \leq 5$,

$$\begin{aligned} e(G) &= e(G_1) + e(G_2) \leq 3|G_1| - 6 + \frac{18(n + 1 - |G_1| - 2)}{7} \\ &= \frac{18(n - 2)}{7} - \frac{(24 - 3|G_1|)}{7} < \frac{18(n - 2)}{7}; \end{aligned}$$

when $|G_1| \geq 6$,

$$e(G) = e(G_1) + e(G_2) \leq \frac{18(|G_1| + |G_2| - 4)}{7} < \frac{18(n - 2)}{7}.$$

So we may assume that G is 2-connected. By Observation 2.1(d) and the fact that $\delta(G) \geq 3$, each face in G is bounded by a cycle.

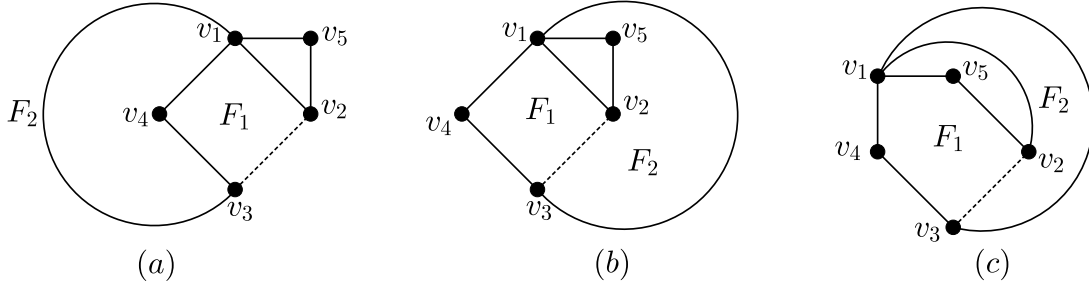


Figure 5: Three possible configurations of H with $v_2v_3 \in E_{4,4} \cup E_{3,5}$.

Assume next that $E_{4,4} \cup E_{3,5} \neq \emptyset$. Let $v_2v_3 \in E_{4,4} \cup E_{3,5}$. Since $\delta(G) \geq 3$, let F_1 and F_2 be the two faces of G having v_2v_3 in common such that the size of F_1 is at least the size of F_2 . Then G must contain a plane subgraph H isomorphic to the graphs depicted in Figure 5(a,b) when $v_2v_3 \in E_{4,4}$, and in Figure 5(c) when $v_2v_3 \in E_{3,5}$, because G is Θ_6 -free. Let H_1 and H_2 be the induced plane subgraphs of G with boundary v_1, v_2, v_5 and v_1, v_3, v_4 , respectively. Then $|H_1| + |H_2| = n + 1$, and $|H_i| \geq 6$ for all $i \in [2]$ because G is Θ_6 -free, 2-connected and $\delta(G) \geq 3$. By the induction hypothesis, $e(H_i) \leq 18(|H_i| - 2)/7$ for all $i \in [2]$. Thus,

$$e(G) = e(H_1) + e(H_2) + |\{v_2v_3\}| < 18(n - 2)/7.$$

We may now further assume that $E_{4,4} \cup E_{3,5} = \emptyset$. Then $e_{4,4} = 0$ and $e_{3,5} = 0$.

It is easy to see that G is not a plane triangulation and so $\sum_{i \geq 4} f_i \geq 0$. We next show that $\sum_{i \geq 5} f_i \neq 0$. Suppose $\sum_{i \geq 5} f_i = 0$. Then $f_3 + f_4 = f$ and $f_4 > 0$. Note that $e_{4,4} = 0$. It follows that every edge of a 4-face of G belongs to $E_{3,4}$, and so G contains a Θ_6 subgraph, a contradiction. Thus $\sum_{i \geq 5} f_i \neq 0$. We may further assume that the outer face of G is neither a 3-face nor a 4-face. Then

$$\begin{aligned} 2e(G) &= 3f_3 + 4f_4 + 5f_5 + \sum_{i \geq 6} if_i \\ &\geq 3f_3 + 4f_4 + 5f_5 + 6(f - f_3 - f_4 - f_5) \\ &= 6f - 3f_3 - 2f_4 - f_5, \end{aligned} \tag{5}$$

which implies that $6f \leq 2e(G) + 3f_3 + 2f_4 + f_5$.

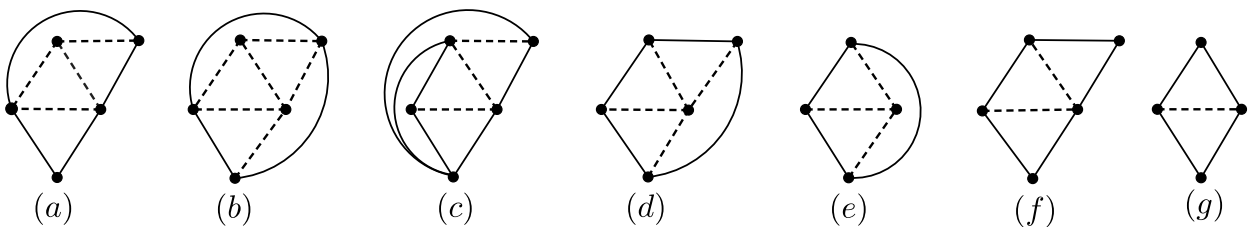


Figure 6: All possible configurations of H_F , where all dashed edges are in $E_{3,3}$, and no solid edges are in $E_{3,3}$.

We next find an upper bound for each of f_3 , f_4 and f_5 . To get an upper bound for f_3 , we first show that $5e_{3,3} \leq 6f_3$. Let F be a 3-face in G with $|E(F) \cap E_{3,3}| \geq 1$. Clearly, $|E(F) \cap E_{3,3}| \leq 3$. Since G is Θ_6 -free and the outer face of G is not a 3-face, there exists a plane subgraph H_F of G with $|H_F| \leq 5$ such that F is a face (not the outer face) of H_F ; all faces of H_F , except the outer face of H_F and any face of H_F that is not a face in G , are 3-faces; and no edges on the boundary of the outer face of H_F and any face of H_F that is not a face in G are in $E_{3,3}$. The possible configurations of H_F are shown in Figure 6. When H_F is isomorphic to the graph depicted in Figure 6(b), H_F contains six edges in $E_{3,3}$ and five 3-faces of G . From all possible configurations of H_F , we see that $e_{3,3} \leq 6f_3/5$. Hence,

$$3f_3 = e_3 + e_{3,3} \leq e_3 + 6f_3/5, \text{ and so } f_3 \leq 5e_3/9. \tag{6}$$

To get an upper bound for f_4 , we next show that $4f_4 \leq 2(e(G) - e_3)$. By Observation 2.1(b,c),

$$4f_4 = e_4 \text{ and } e(G) \geq e_3 + e_4 - e_{3,4}. \tag{7}$$

We next show that $e_{3,4} \leq e(G) - e_3$.

This is trivially true when $e_{3,4} = 0$. Assume that $e_{3,4} \neq 0$. Let F and F' be a 4-face and a 3-face in G , respectively, such that F and F' share an edge in common. We may assume that F

has vertices v_1, v_2, v_3, v_4 in order and F' has vertices v_1, v_4, v_5 in order. Note that F and F' are not outer face in G . Observe that if $v_i v_{i+1}$ belongs to $E_{3,4}$ for any $i \in \{1, 2, 3\}$, then $v_i v_{i+1}$ belongs to the 4-face F and the 3-face with vertices v_i, v_{i+1}, v_5 in order, else G would not be Θ_6 -free. Since $n \geq 7$, there exists some $k \in \{1, 2, 3\}$ such that $v_k v_{k+1} \notin E_{3,4}$. Then $v_k v_{k+1} \in E_{4,j}$ for some $j \geq 5$ because $e_{4,4} = 0$. We next show that F has at most two edges in $E_{3,4}$. Suppose $|E(F) \cap E_{3,4}| = 3$. We may assume that $k = 2$. Then $v_1 v_2, v_3 v_4 \in E_{3,4}$. Thus $v_1 v_2$ belongs to the 4-face F and the 3-face with vertices v_1, v_2, v_5 in order; and $v_3 v_4$ belongs to the 4-face F and the 3-face with vertices v_3, v_4, v_5 in order. Since G is Θ_6 -free, we see that $v_5 v_2 \in E_{3,j}$ for some $j \geq 6$, $v_5 v_3 \in E_{3,j}$ for some $j \geq 6$, and $v_2 v_3 \in E_{4,j}$ for some $j \geq 6$. But then $G + v_2 v_4$ is Θ_6 -free, contrary to the choice of G . Thus F has at most two edges in $E_{3,4}$, and so F has at least two edges not in E_3 . This holds for each 4-face in G . Hence, $e_{3,4} \leq e(G) - e_3$.

By (7),

$$4f_4 = e_4 \leq e(G) - e_3 + e_{3,4} \leq 2(e(G) - e_3). \quad (8)$$

Note that $e_{3,5} = 0$. By Observation 2.1(a), $e_{5,5} \leq e_5 \leq e(G) - e_3$. By Observation 2.1(b),

$$5f_5 = e_5 + e_{5,5} \leq 2(e(G) - e_3). \quad (9)$$

Combining $e_3 \leq e(G)$ with the upper bounds on f_3, f_4, f_5 given in (6), (8), (9), we have

$$\begin{aligned} 6f &\leq 2e(G) + 3f_3 + 2f_4 + f_5 \\ &\leq 2e(G) + 5e_3/3 + (e(G) - e_3) + 2(e(G) - e_3)/5 \\ &= 17e(G)/5 + 4e_3/15 \\ &\leq 11e(G)/3. \end{aligned}$$

It follows that $f \leq 11e(G)/18$. By Euler's formula, $n - 2 \geq e(G) - f \geq 7e(G)/18$. Hence $e(G) \leq 18(n - 2)/7$, as desired. \blacksquare

Corollary 2.5 *Let K_5^- be the graph obtained from K_5 by deleting one edge. Then*

(a) $ex_{\mathcal{P}}(n, \Theta_6 \cup \{K_5^-\}) \leq 12(n - 2)/5$ for all $n \geq 7$.

(b) $ex_{\mathcal{P}}(n, \Theta_6) < 18(n - 2)/7$ for all $n \geq 10$.

Proof. To prove (a), let G be an extremal plane graph for $n \geq 7$ and $\Theta_6 \cup \{K_5^-\}$. We prove that $e(G) \leq 12(n - 2)/5$ by induction on n . Since G is Θ_6 -free, by Theorem 2.4, $e(G) \leq \lfloor 18(n - 2)/7 \rfloor = 12(n - 2)/5$ when $n = 7$. So we may assume that $n \geq 8$. Similar to the proof of Theorem 2.4, we see that $e_{3,3} \leq f_3$, because G is K_5^- -free and so no H_F is isomorphic to the graphs depicted in

Figure 6(b, c). By Observation 2.1(b), $f_3 \leq e_3/2$. This, together with the upper bounds for f_4, f_5 given in (8) and (9), implies that

$$\begin{aligned}
6f &\leq 2e(G) + 3f_3 + 2f_4 + f_5 \\
&\leq 2e(G) + 3e_3/2 + e(G) - e_3 + 2(e(G) - e_3)/5 \\
&= 17e(G)/5 + e_3/10 \\
&\leq 7e(G)/2.
\end{aligned}$$

It follows that $f \leq 7e(G)/12$. By Euler's formula, $n - 2 \geq e(G) - f \geq 5e(G)/12$. Hence $e(G) \leq 12(n - 2)/5$.

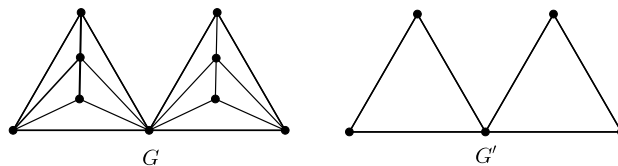


Figure 7: An example of constructing the graph G' from a graph G .

To prove (b), let G be an extremal plane graph for Θ_6 and $n \geq 10$. By Theorem 2.4, $e(G) \leq 18(n - 2)/7$. Suppose $e(G) = 18(n - 2)/7$. By Corollary 2.5(a), G is not K_5^- -free. From the proof of Theorem 2.4, we see that equality in $e(G) \leq 18(n - 2)/7$ is achieved for n if and only if all the equalities hold in (5), (6), (8), (9) and in $e_3 \leq e(G)$. This implies that $e(G) = 18(n - 2)/7$ for n if and only if G is a 2-connected Θ_6 -free plane graph on n vertices satisfying: G consists entirely of K_5^- 's and 6-faces, no two K_5^- 's share an edge, and no two 6-faces have an edge in common. Let G' be the graph obtained from G by deleting the two vertices not on the outer face in each $H_F = K_5^-$, an example is shown in Figure 7. Then G' consists of 3-faces and 6-faces such that each edge of G' belongs to one 3-face and one 6-face. Clearly, G' is 2-connected because G is 2-connected. Let f'_i be the number of i -faces in G' . Let $f' = \sum_{i \geq 1} f'_i$. Then $3f'_3 = e(G') = 6f'_6$ and $f' = f'_3 + f'_6$. Thus $|G'| - 2 = e(G') - f' = e(G')/2$ and so $e(G') = 2|G'| - 4$, which implies that $\delta(G') \leq 3$. Since G' is 2-connected, we have $\delta(G') \geq 2$. Note that each vertex of G' must have even degree because the adjacent faces are alternatively of size 3 and size 6. Thus $\delta(G') = 2$. Let $v \in V(G')$ be a vertex of degree two in G' . Let u_1vu_2 and $u_1vu_2u_3u_4u_5$ be the vertices in order on the boundary of the two adjacent faces containing v , respectively. Then $G'[\{u_1, v, u_2, u_3, u_4, u_5\}]$ contains a graph in Θ_6 as a subgraph. Thus G is not Θ_6 -free, a contradiction. \blacksquare

It is worth noting that every C_6 -free graph is certainly Θ_6 -free. Hence, $ex_{\mathcal{P}}(n, C_6) \leq ex_{\mathcal{P}}(n, \Theta_6)$. Corollary 2.6 follows immediately from Theorem 2.4.

Corollary 2.6 $ex_{\mathcal{P}}(n, C_6) \leq 18(n - 2)/7$ for all $n \geq 6$, with equality when $n = 9$.

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